

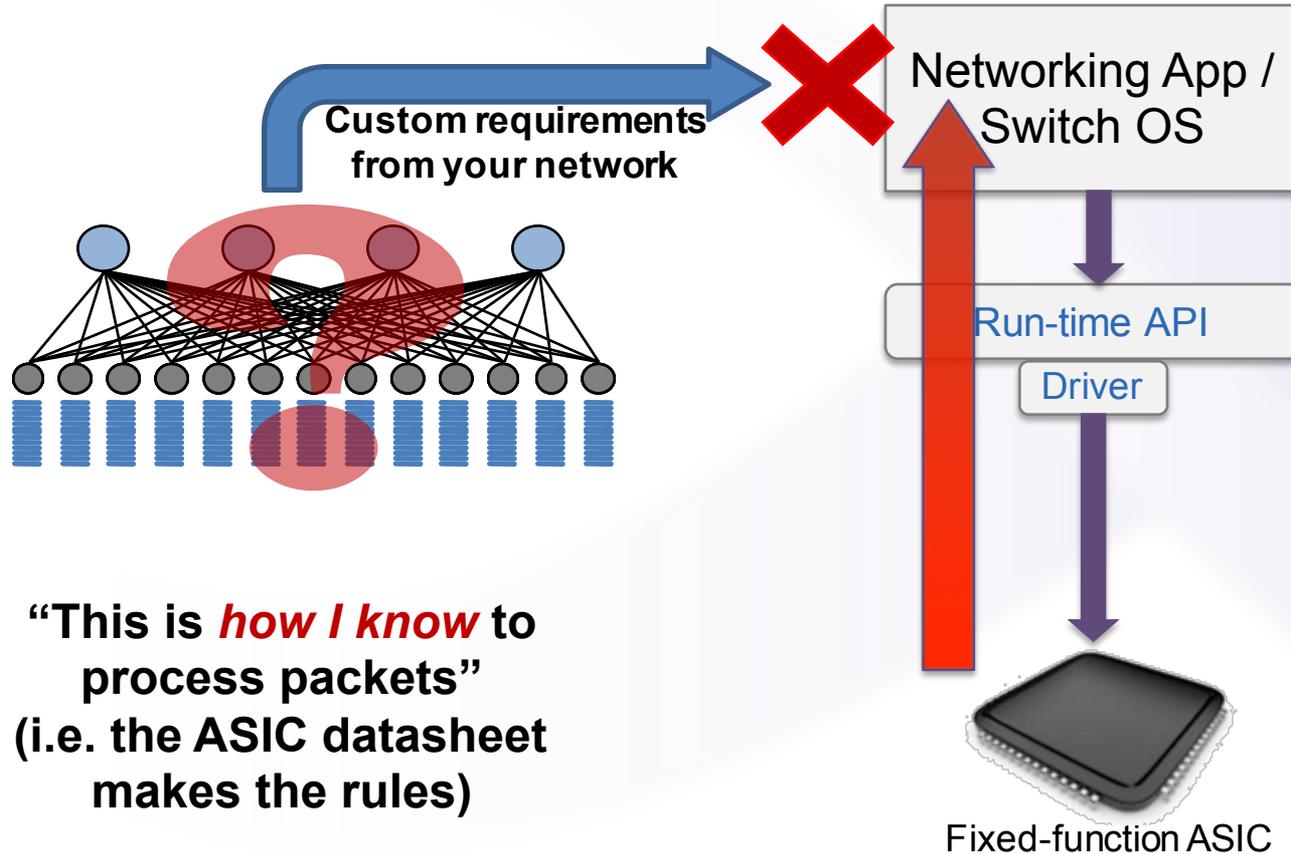
Programming The Network Data Plane in P4

**SIGCOMM'16
Tutorial**

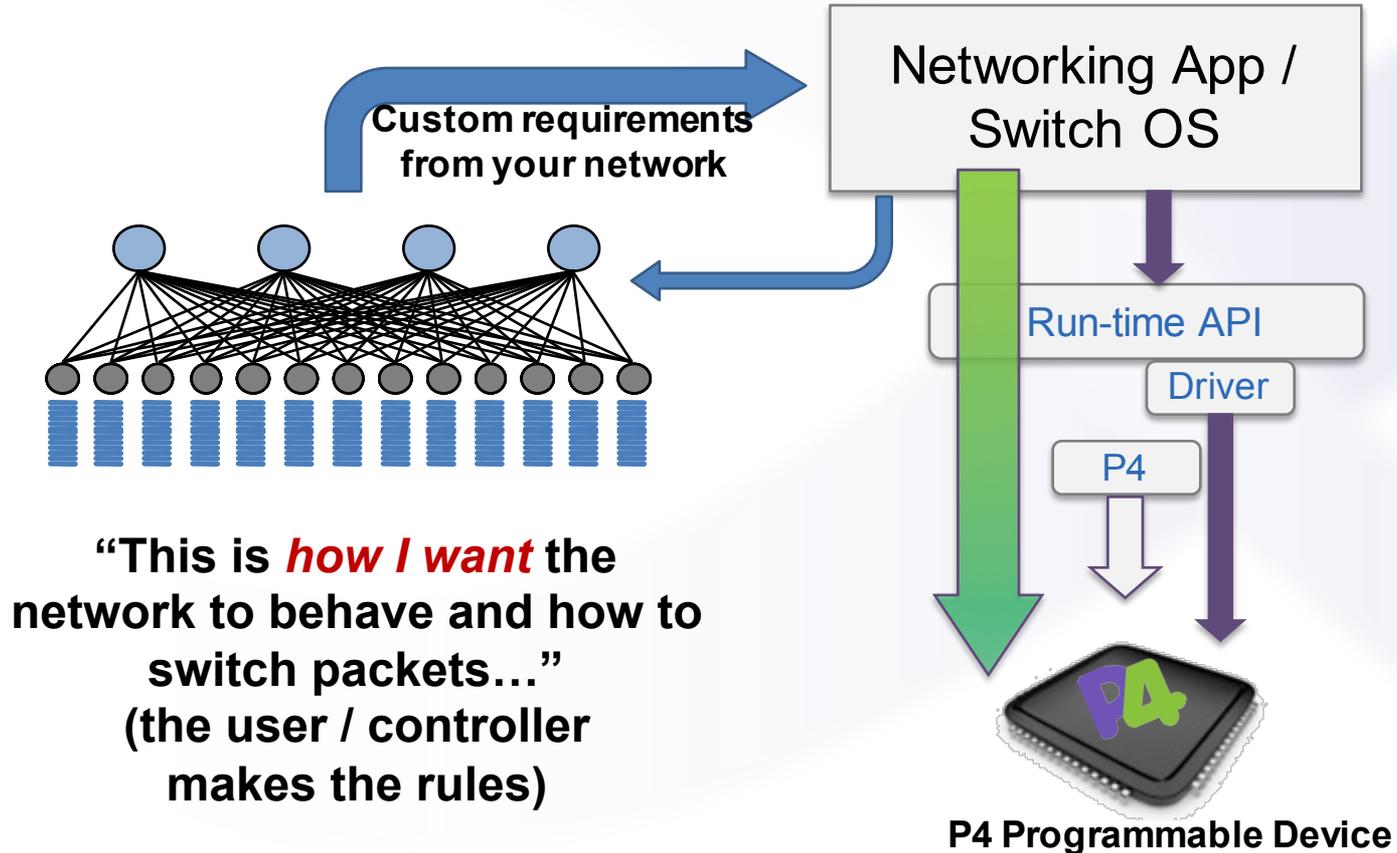
**Barefoot Networks
Aug 2016**

P4 Introduction

Status Quo: Bottom-up design



A Better Approach: Top-down design



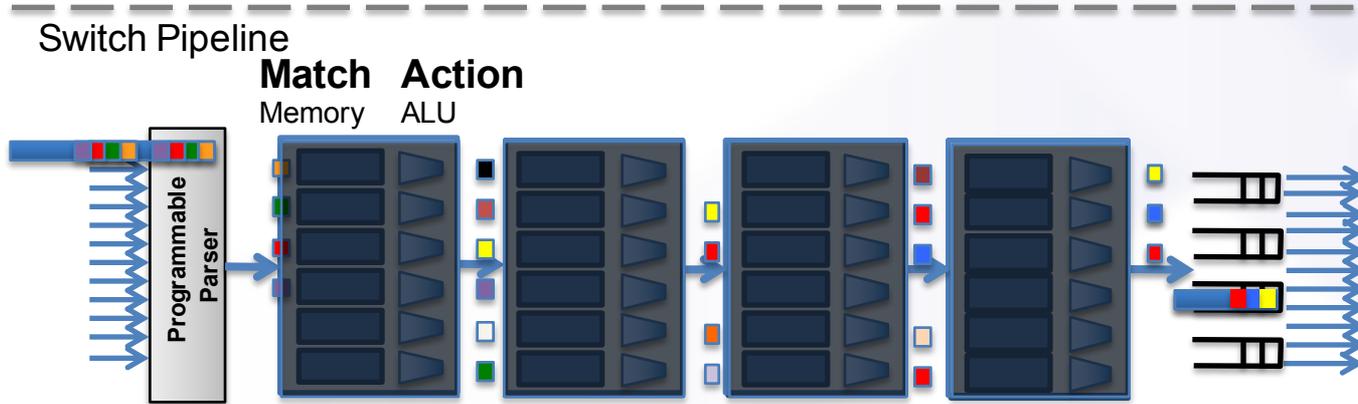
Programmable Network Devices

- **PISA: Flexible Match+Action ASICs**
 - Intel Flexpipe, Cisco Doppler, Cavium (Xpliant), Barefoot Tofino, ...
- **NPU**
 - EZchip, Netronome, ...
- **CPU**
 - Open Vswitch, eBPF, DPDK, VPP...
- **FPGA**
 - Xilinx, Altera, ...

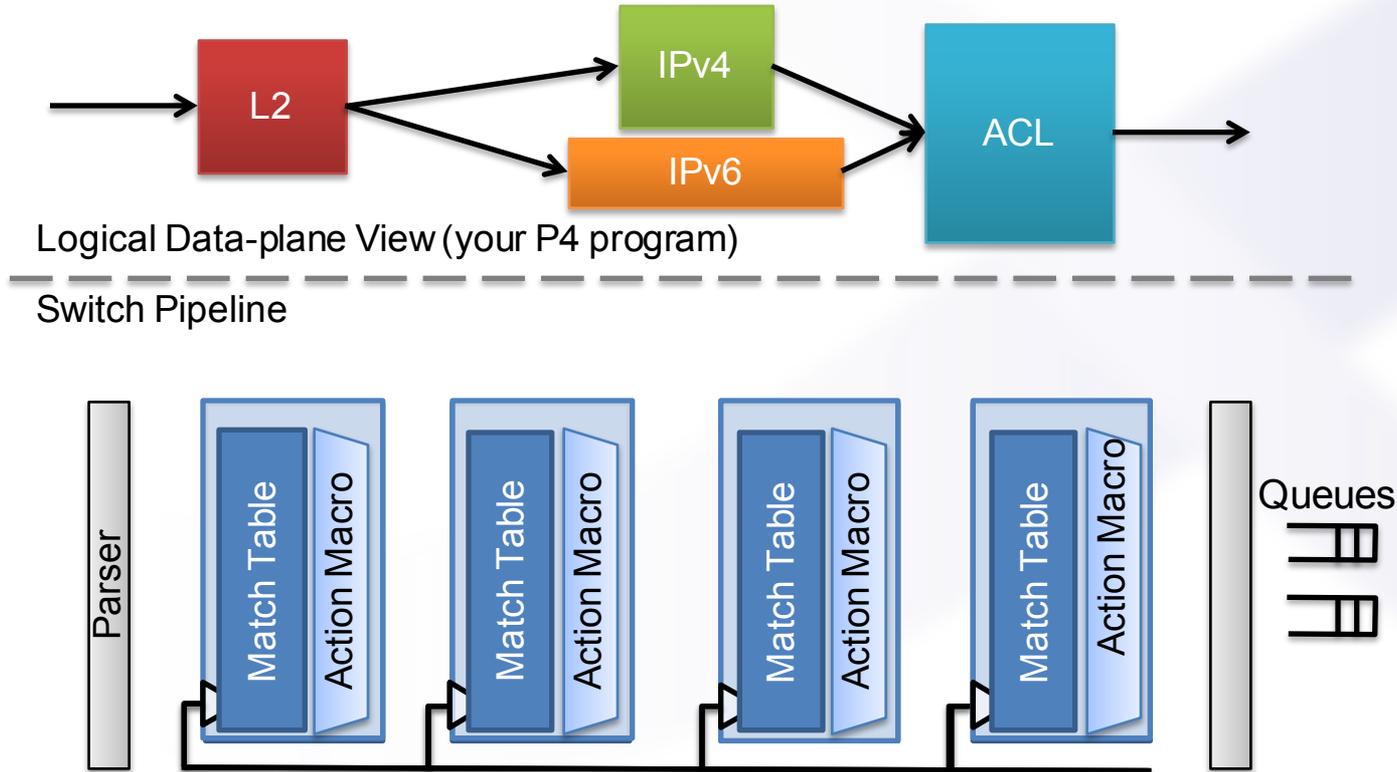
These devices let us tell them how to process packets.

Why we call it Protocol Independent Packet Processing

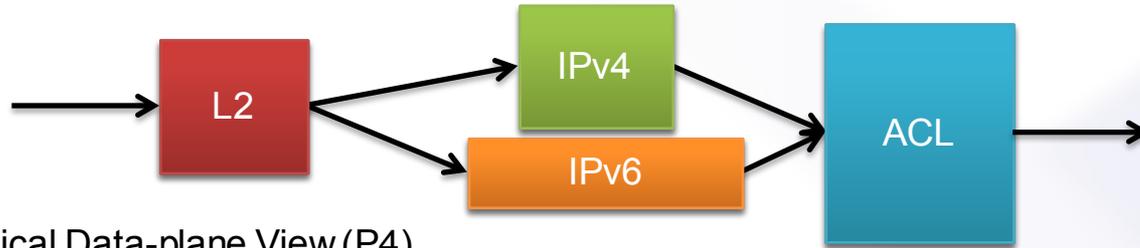
Protocol-Independent Switch Architecture (PISA)



Protocol-Independent Switch Architecture (PISA)

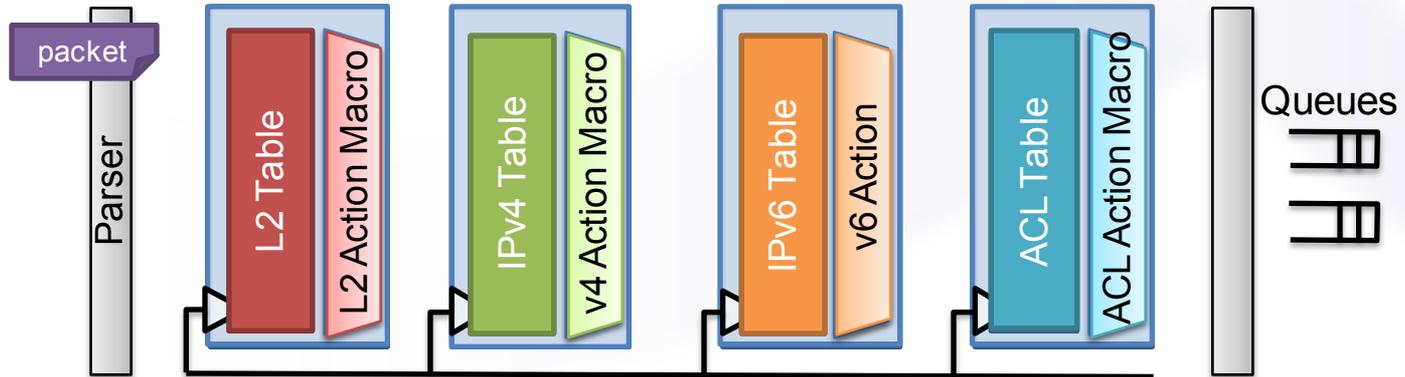


Mapping to Physical Resources

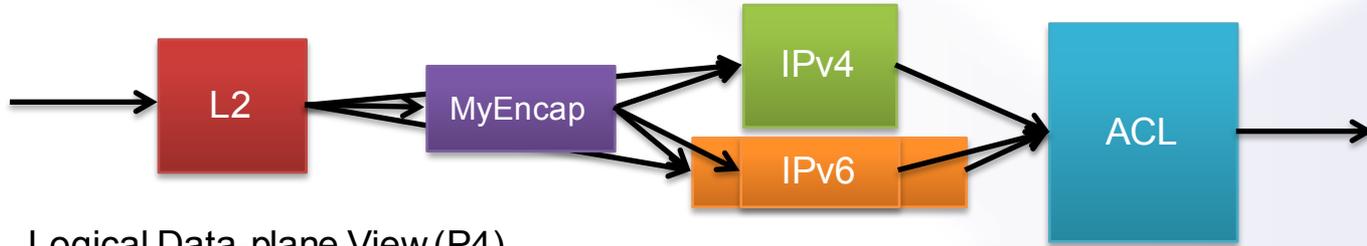


Logical Data-plane View (P4)

Switch Pipeline

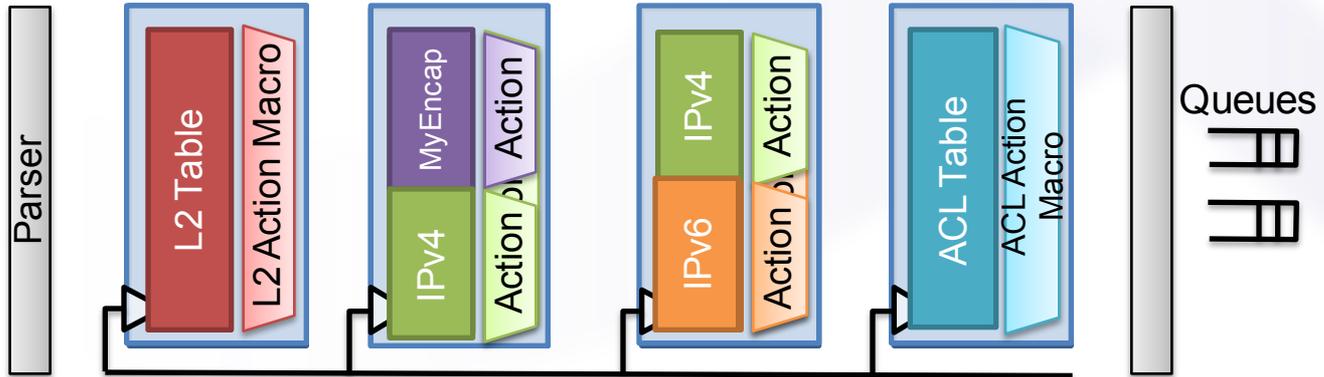


Re-configurability



Logical Data-plane View (P4)

Switch Pipeline



P4: Three Goals

Protocol independence

- Define a packet parser
- Define a set of typed match+action tables

Target independence

- Program without knowledge of packet-processing device details
- Let compilers configure the target device

In-field Re-configurability

- Allow the users to change parsing and processing program in the field

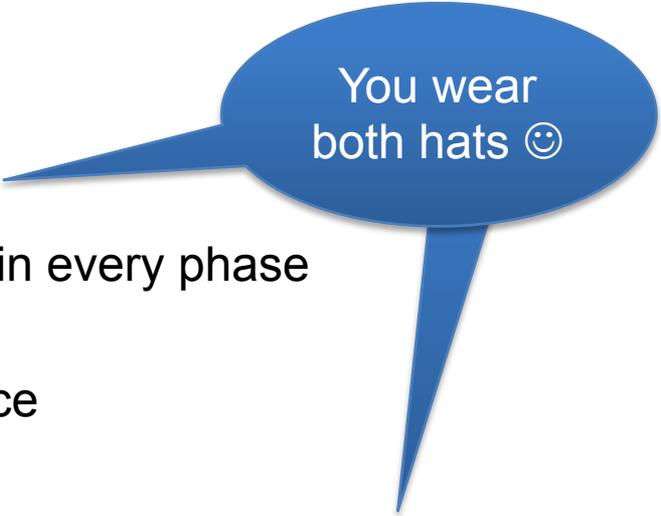
What does this mean?

■ To network device vendors

- S/W programming practices and tools used in every phase
- Extremely fast iteration and feature release
- Differentiation in capabilities and performance
- Can fix even data-plane bugs in the field

■ To large on-line service providers and carriers

- No more “black boxes” in the “white boxes”
- Your devs can program, test, and debug your network devices all the way down
- You keep your own ideas



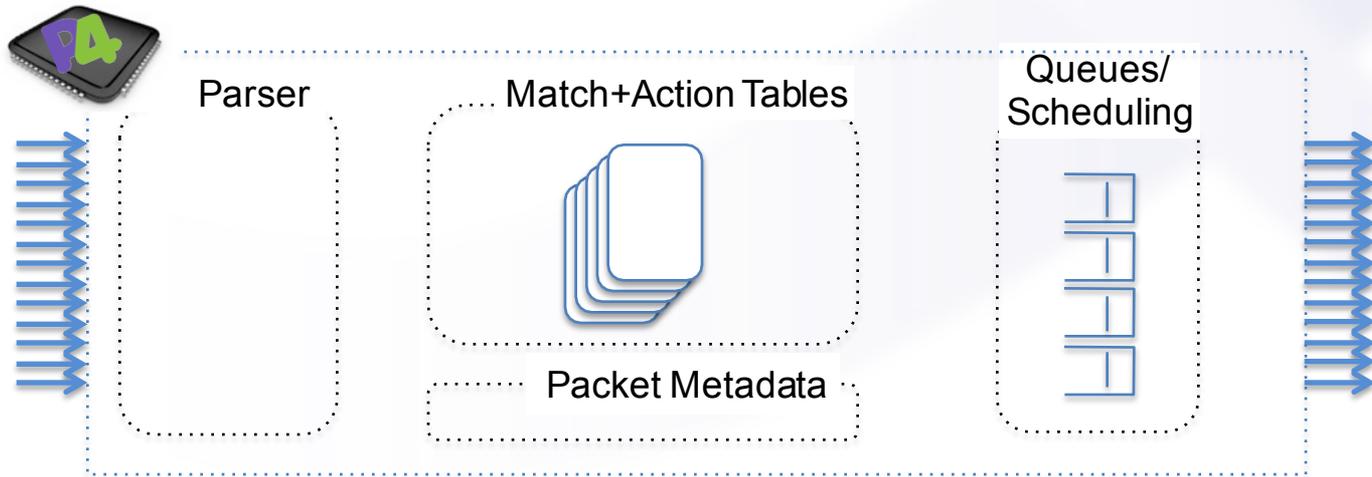
You wear
both hats 😊

Key benefits of programmable forwarding

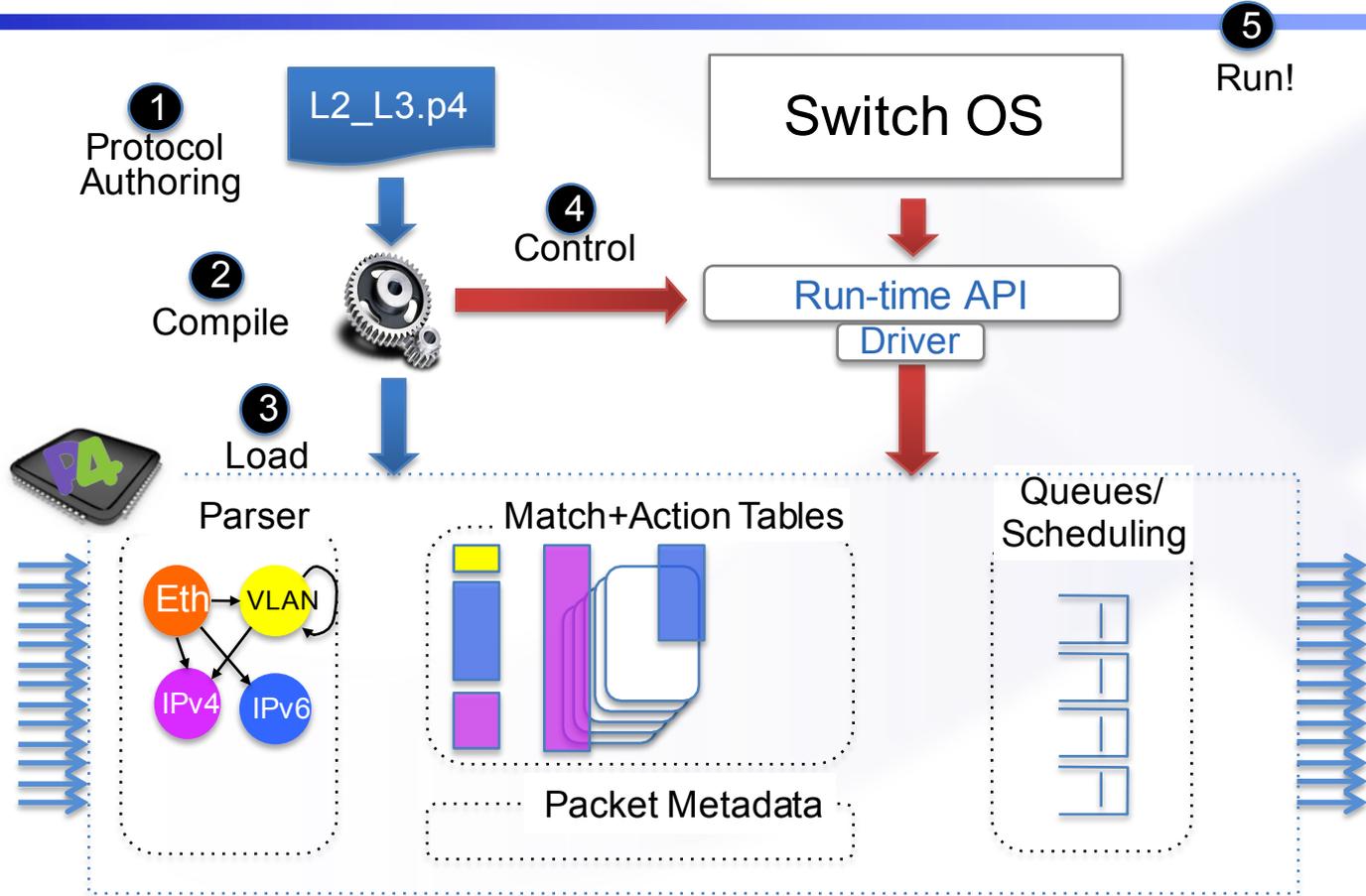
1. **New features**: Realize new protocols and behaviors very quickly
2. **Reduce complexity**: Remove unnecessary features and tables
3. **Efficient use of H/W resources**: Achieve biggest bang for buck
4. **Greater visibility**: New diagnostics, telemetry, OAM, etc.
5. **Modularity**: Compose forwarding behavior from libraries
6. **Portability**: Specify forwarding behavior once; compile to many devices
7. **Own your own network**: No need to wait for next chips or systems

P4-Based Workflow

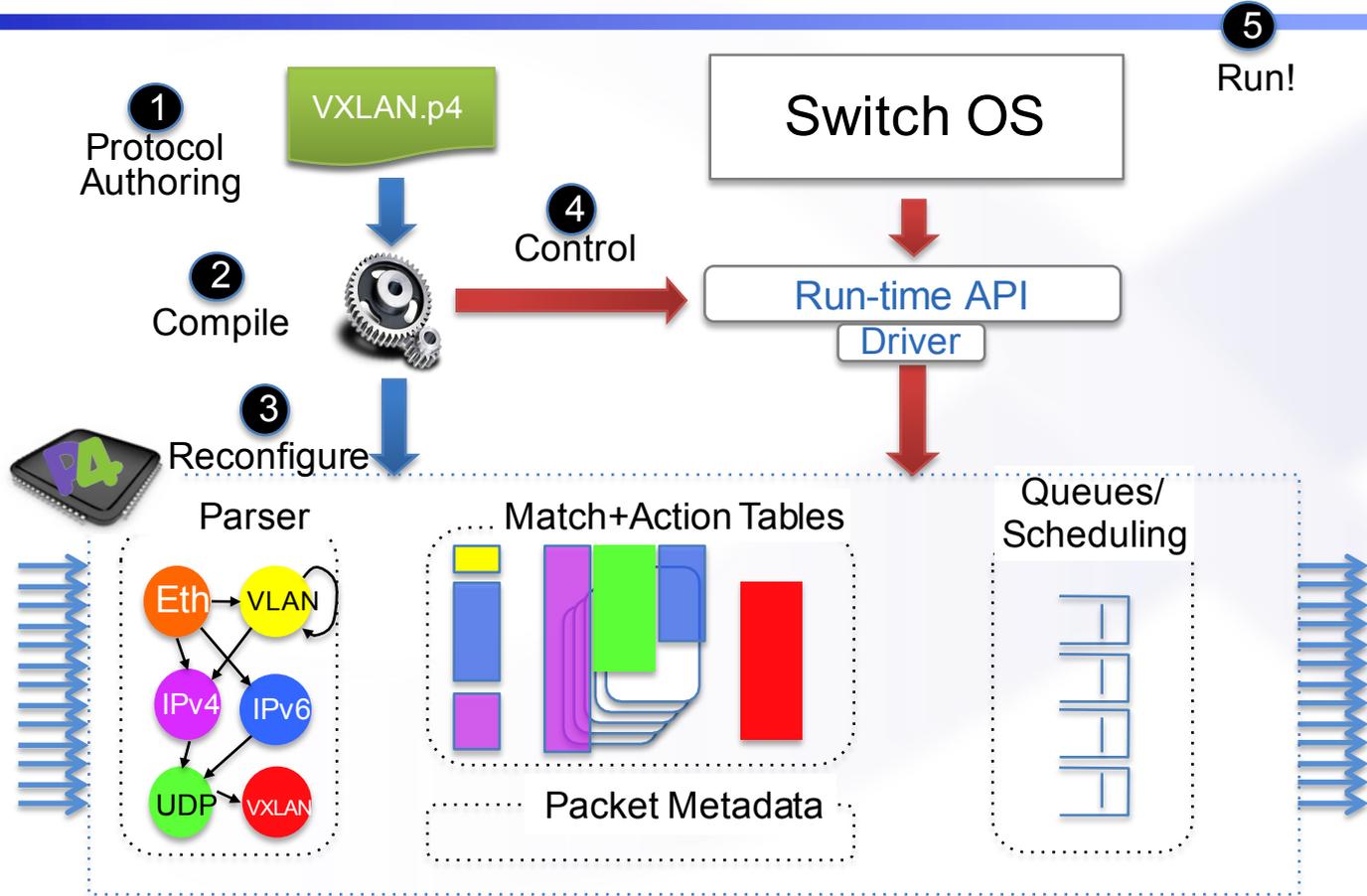
- **Device is not yet programmed**
 - Does not know about any packet formats or protocols

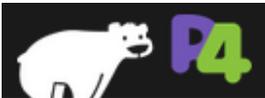


P4-Based Workflow



P4-Based Workflow





The P4 Language Consortium

- Consortium of academic and industry members
- Open source, evolving, domain-specific language
- Permissive Apache license, code on GitHub today
- Membership is free: contributions are welcome
- Independent, set up as a California nonprofit

Protocol Independent
P4 programs specify how a switch processes packets.

Target Independent
P4 is suitable for describing everything from high-performance forwarding ASICs to software switches.

Field Reconfigurable
P4 allows network engineers to change the way their switches process packets after they are deployed.

```
table routing {
  reads {
    ipv4.dstAddr : lpm;
  }
  actions {
    do_drop;
    route_ipv4;
  }
  size: 2048;
}

control ingress {
  apply(routing);
}
```

 **TRY IT** Get the code from P4factory



P4.org Membership



Original P4 Paper Authors:



Operators



Systems



Targets



Academia

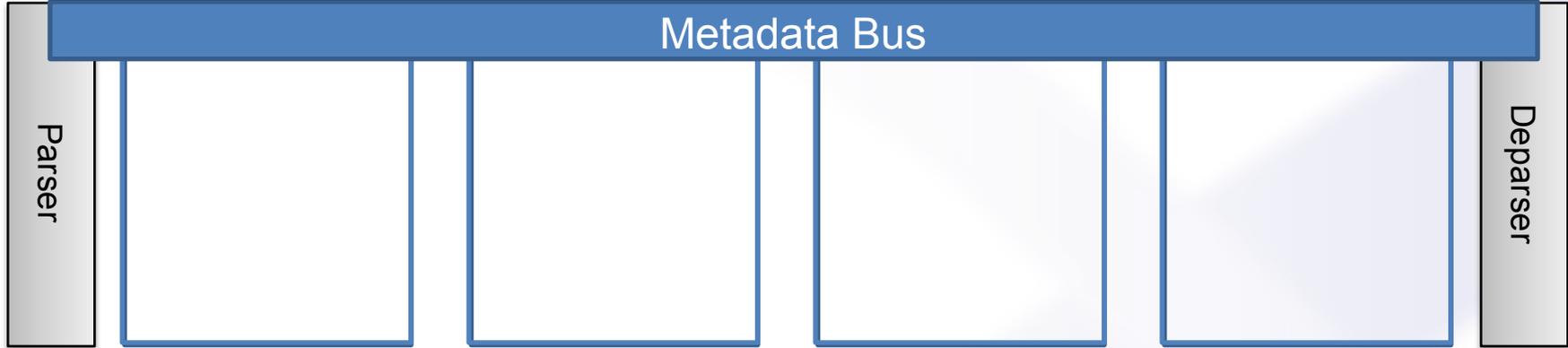


- **Open source**, evolving, domain-specific language
- Permissive Apache license, code on GitHub today
- **Membership is free**: contributions are welcome
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P4 Concepts

- **Pipeline**
 - Parser / Deparser
 - Match-Action Tables

The anatomy of a basic pipeline

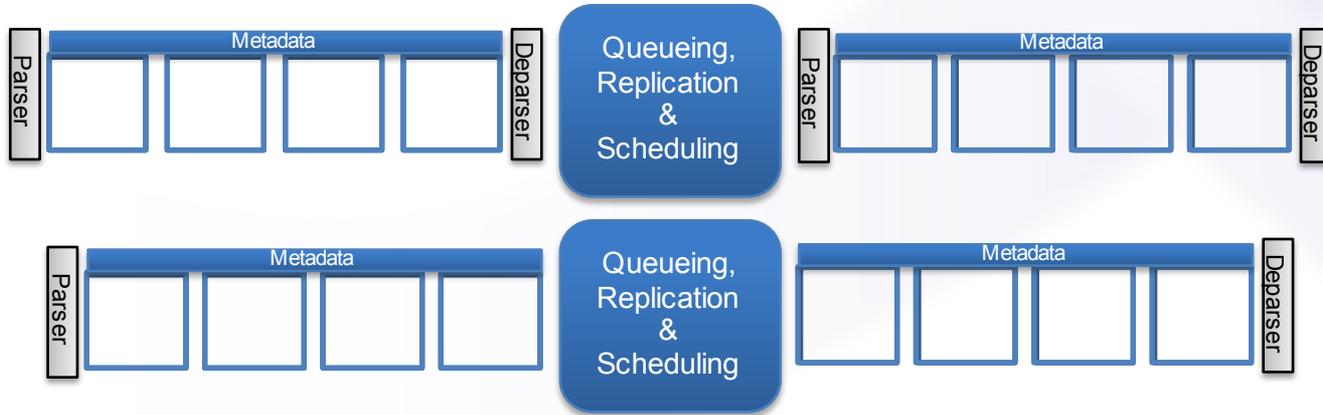


- **Parser**
 - Converts packet data into a metadata (Parsed Representation)
- **Match+Action Tables**
 - Operate on metadata
- **Deparser**
 - Converts metadata back into a serialized packet
- **Metadata Bus**
 - Carries the information within the pipeline

All are optional

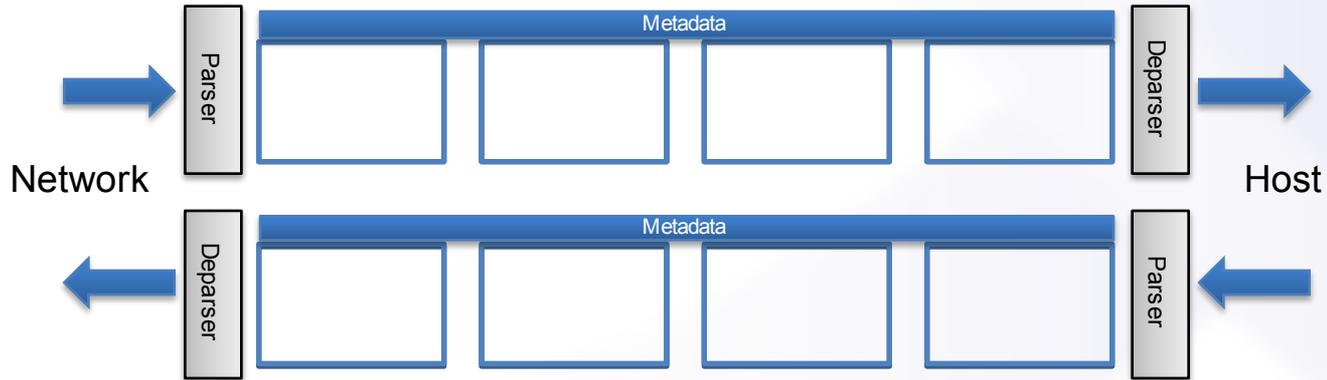
Anatomy of a Switch

- **Ingress Pipeline**
- **Egress Pipeline**
- **Traffic Manager**
 - N:1 Relationships: Queueing, Congestion Control
 - 1:N Relationships: Replication
 - Scheduling



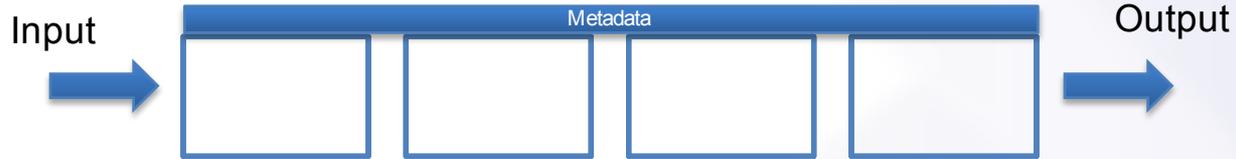
Anatomy of a NIC

- Single or Dual Pipeline



Anatomy of Protocol Plugin

- **Single, “Bare” Pipeline**
 - No parsing/deparsing, just processing



P4 Program Sections

program.p4

Data Declarations

```
header_type ethernet_t { ... }
header_type l2_metadata_t { ... }

header ethernet_t ethernet;
header vlan_tag_t vlan_tag[2];
metadata l2_metadata_t l2_meta;
```

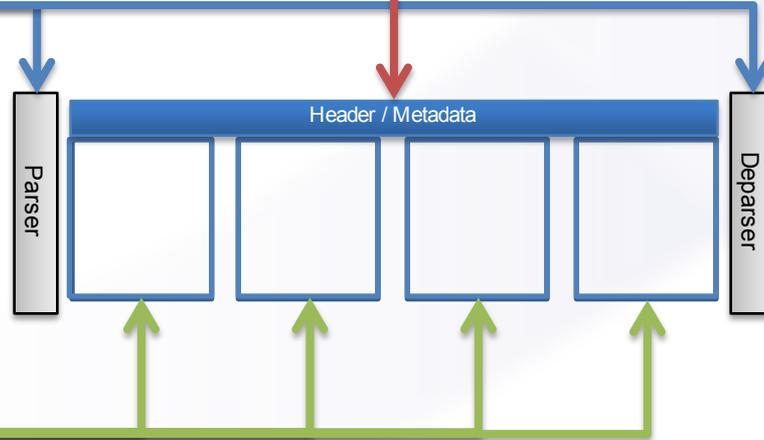
Parser Program

```
parser parse_ethernet {
  extract(ethernet);
  return switch(ethernet.ethertype) {
    0x8100 : parse_vlan_tag;
    0x0800 : parse_ipv4;
    0x8847 : parse_mpls;
    default: ingress;
  }
}
```

Table + Control Flow Program

```
table port_table { ... }

control ingress {
  apply(port_table);
  if (l2_meta.vlan_tags == 0) {
    process_assign_vlan();
  }
}
```



P4 program defines what each table CAN do

Control Plane Roles

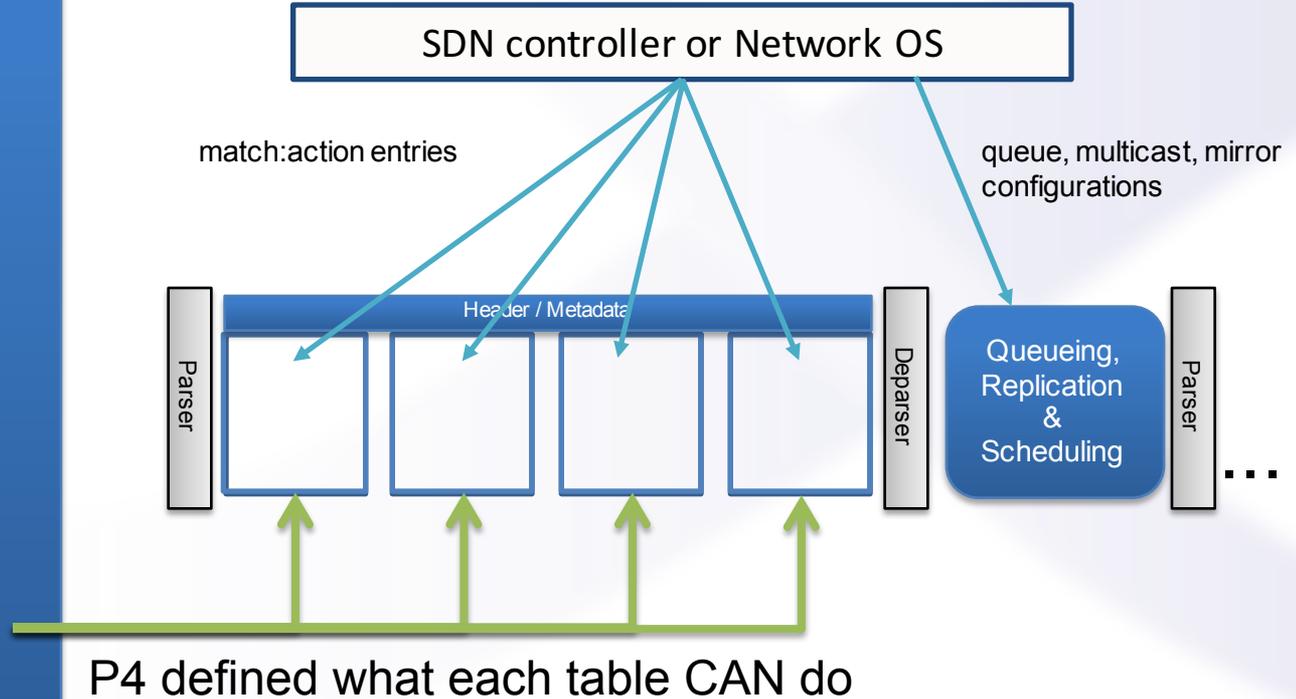
program.p4

Data Declarations

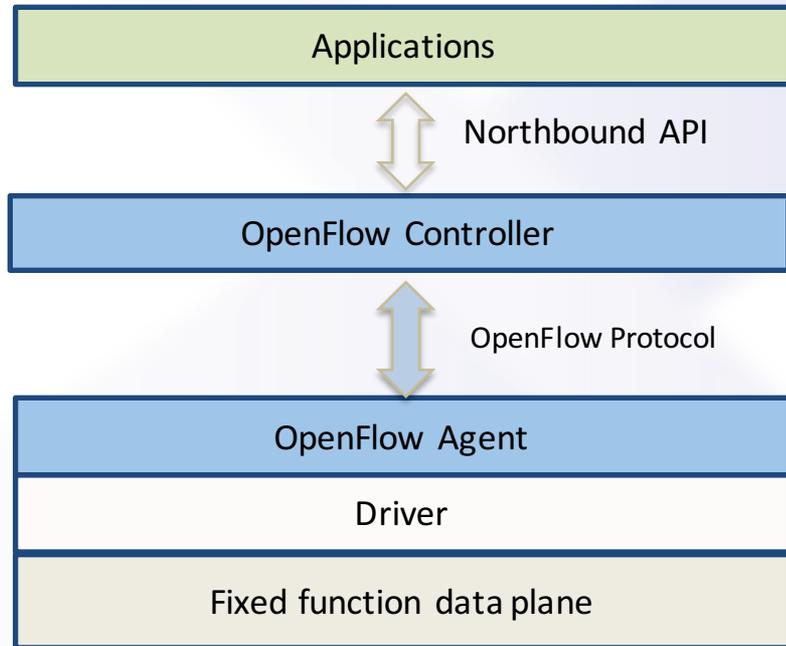
Parser Program

Table + Control Flow Program

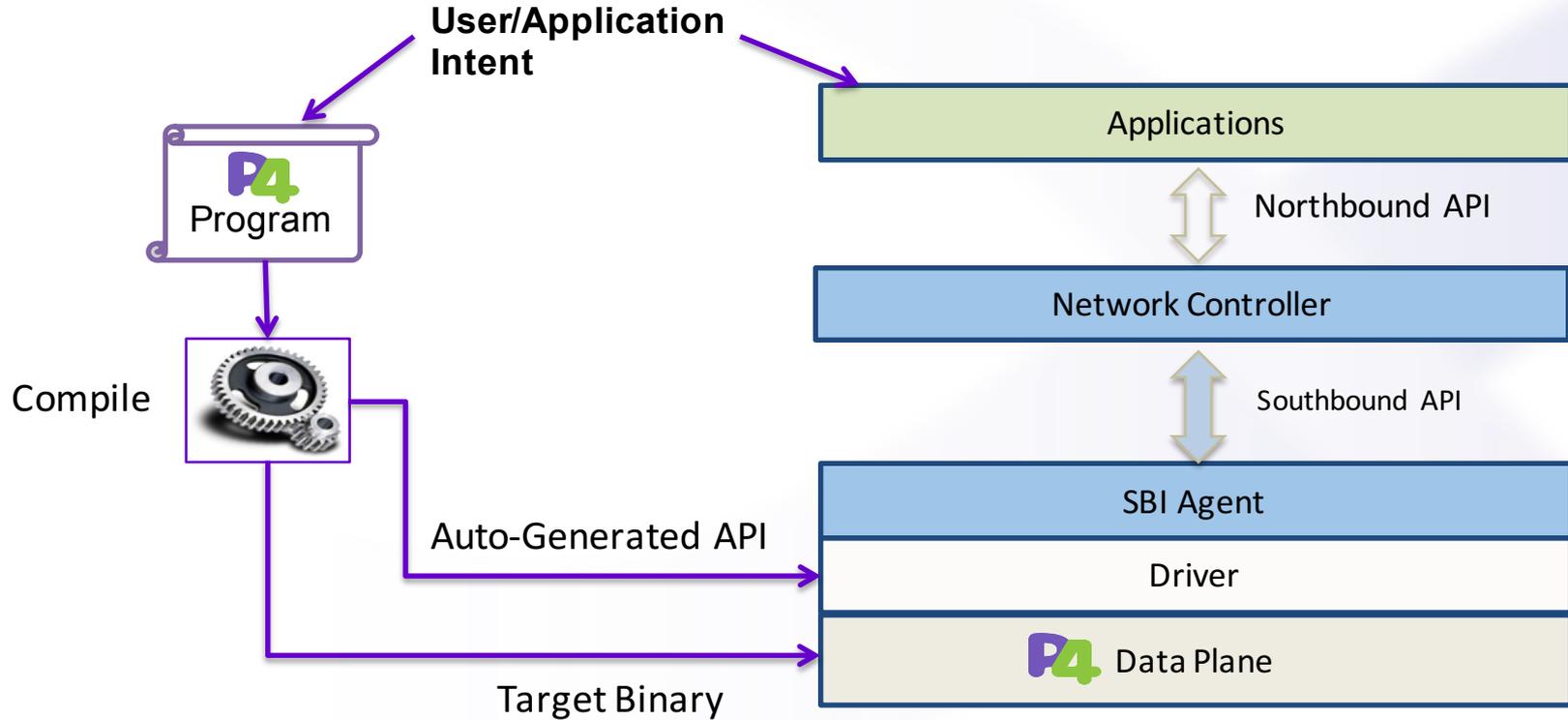
Control plane or NOS decides **switch runtime behavior**



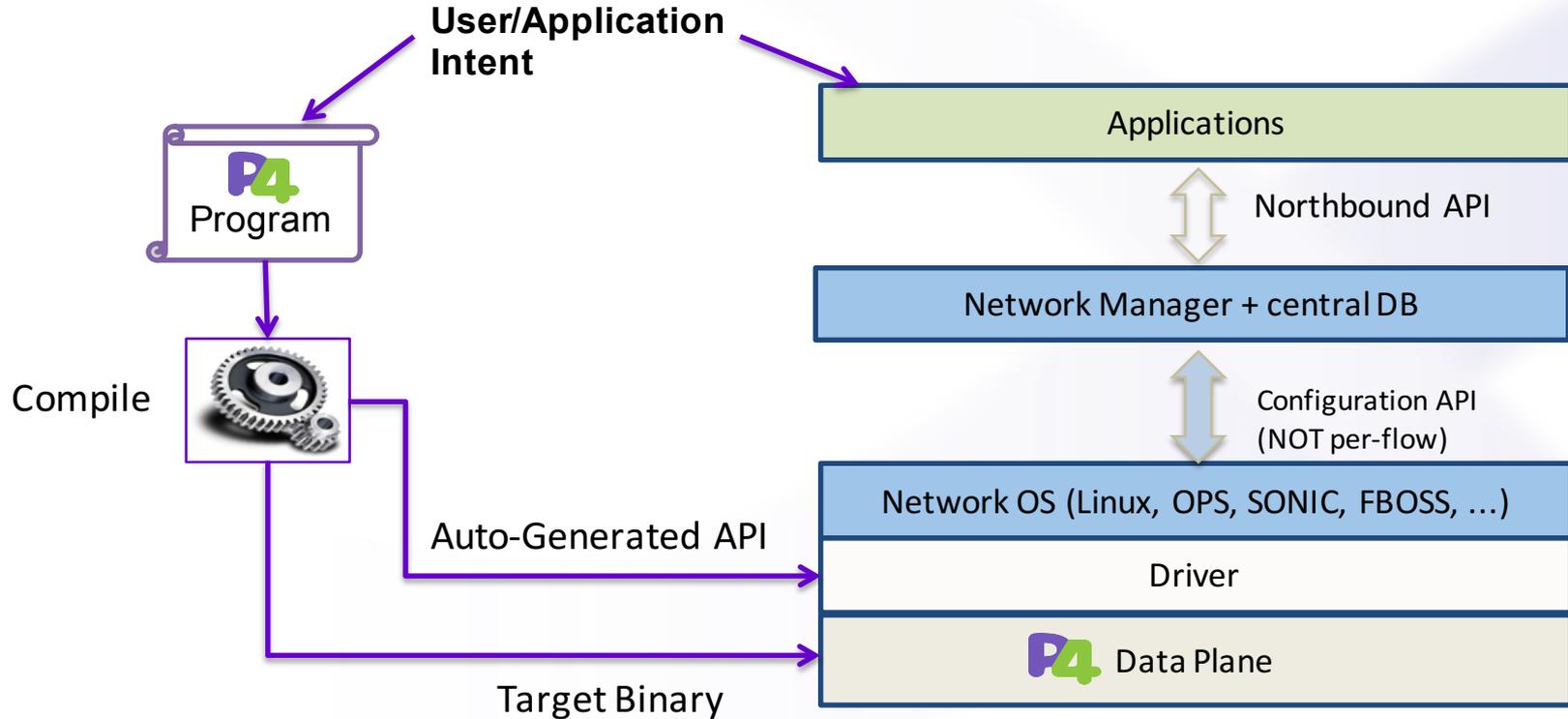
P4 & OpenFlow: Traditional SDN before P4



P4 with OpenFlow



P4 with Network OS



P4 Constructs

- P4 Spec v1.0.2+, v1.1.0-

P4 Language Components

- **Data declarations**

- Packet Headers and Metadata

- **Parser Programming**

- Parser Functions (Parser states)
- Checksum Units

- **Packet Flow Programming**

- Actions
 - Primitive and compound actions
 - Counters, Meters, Registers
- Tables
 - Match keys
 - Attributes
- Control Functions (Imperative Programs)

No: pointers, loops, recursion, floating point

Headers and Fields (Packet)

Example: Declaring packet headers

```
header_type ethernet_t {  
    fields {  
        dstAddr    : 48;  
        srcAddr    : 48;  
        etherType  : 16;  
    }  
}
```

Header Type
Declarations

```
header_type vlan_tag_t {  
    fields {  
        pcp        : 3;  
        cfi        : 1;  
        vid        : 12;  
        etherType  : 16;  
    }  
}
```

Actual Header
Instantiation

```
header ethernet_t ethernet;  
header vlan_tag_t vlan_tag[3];
```

Handy Arrays for
Header Stacks

Headers and Fields (Metadata)

Example: Declaring Metadata

```
header_type ingress_metadata_t {  
    fields {  
        /* Inputs */  
        ingress_port      : 9; /* Available prior to parsing */  
        packet_length     : 16; /* Might not be always available */  
        instance_type     : 2; /* Normal, clone, recirculated */  
        ingress_global_tstamp : 48;  
        parser_status     : 8; /* Parsing Error */  
  
        /* Outputs from Ingress Pipeline */  
        egress_spec       : 16;  
        queue_id          : 9;  
    }  
}  
  
metadata ingress_metadata_t ingress_metadata;
```

Metadata is a header
too

Actual Metadata
Instantiation

Metadata vs. Packet Headers

- **Layout definition**

- Packet header declarations define both the fields and the actual layout in the packet.
- Layout is not defined for metadata

- **Byte Alignment**

- Packet header length must be a multiple of 8 bits
- No special requirements for metadata

- **Validity**

- Packet headers are valid only if present in the packet
- Metadata is ALWAYS valid
 - Default value is either 0 or can be specified explicitly

- **Acceptable fields**

- Packet headers can contain calculated and variable length fields

Variable-Length Fields

Example: Declaring IPv4 packet header

```
header_type ipv4_t {
    fields {
        version      : 4;
        ihl          : 4;
        diffserv     : 8;
        totalLen     : 16;
        identification : 16;
        flags        : 3;
        fragOffset   : 13;
        ttl          : 8;
        protocol     : 8;
        hdrChecksum  : 16;
        srcAddr      : 32;
        dstAddr      : 32;
        options      : *;
    }
    length      : (ihl << 2);
    max_length  : 60;
}
```

Variable-length Field

Calculated, based on
another field

Defining a Parser Tree

Example: Simple Parser for L2/L3 Packets

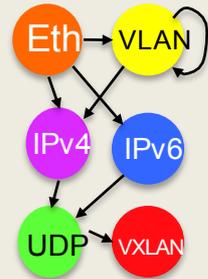
```
header ethernet_t ethernet;  
header vlan_tag_t vlan_tag[2];  
header ipv4_t ipv4;  
header ipv6_t ipv6;
```

```
parser start {  
    extract(ethernet);  
    return select(latest.etherType) {  
        0x8100, 0x9100 : parse_vlan_tag;  
        0x0800       : parse_ipv4;  
        0x86DD       : parse_ipv6;  
        default      : ingress;  
    }  
}
```

Transitions to the next parser states. Prioritized by order

```
parser parse_vlan_tag {  
    extract(vlan_tag[next]);  
    return select(latest.etherType) {  
        0x8100 mask 0xEFFF : parse_vlan_tag;  
        0x0800           : parse_ipv4;  
        0x86DD           : parse_ipv6;  
        default          : ingress;  
    }  
}
```

This is not a reserved word, but a name of the Control Flow Function



Defining a Parser Tree (cont.)

Example: Simple Parser for L2/L3 Packets

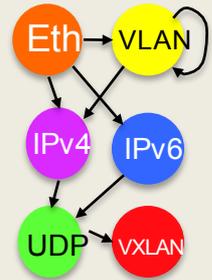
```
header ethernet_t ethernet;  
header vlan_tag_t vlan_tag[2];  
header ipv4_t ipv4;  
header ipv6_t ipv6;
```

```
parser start {  
    extract(ethernet);  
    return select(latest.etherType) {  
        0x8100, 0x9100 : parse_vlan_tag;  
        0x0800         : parse_ipv4;  
        0x86DD         : parse_ipv6;  
        default       : ingress;  
    }  
}
```

```
parser parse_vlan_tag {  
    extract(vlan_tag[next]);  
    return select(latest.etherType) {  
        0x8100 mask 0xEFFF : parse_vlan_tag;  
        0x0800             : parse_ipv4;  
        0x86DD             : parse_ipv6;  
        default           : ingress;  
    }  
}
```

```
parser parse_ipv4 {  
    extract(ipv4);  
    return ingress;  
}
```

```
parser parse_ipv6 {  
    extract(ipv6);  
    return ingress;  
}
```



Using Calculated Fields

Example: Calculated fields for IPv4

```
field_list ipv4_checksum_list {
    ipv4.version;
    ipv4.ihl;
    ipv4.diffserv;
    ipv4.totalLen;
    ipv4.identification;
    ipv4.flags;
    ipv4.fragOffset;
    ipv4.ttl;
    ipv4.protocol;
    ipv4.srcAddr;
    ipv4.dstAddr;
}

field_list_calculation ipv4_checksum {
    input      { ipv4_checksum_list; }
    algorithm  : csum16;
    output_width : 16;
}

calculated_field ipv4_hdrChecksum {
    verify ipv4_checksum;
    update ipv4_checksum;
}
```

```
parser parse_ipv4 {
    extract(ipv4);
    return ingress;
}

parser_exception p4_pe_checksum {
    return parser_drop;
}
```



Predefined parser
state

Multi-field select statement

Example: Ipv4 Header Parsing

```
parser parse_ipv4 {
  extract(ipv4);
  set_metadata(ipv4_metadata.lkp_ipv4_sa, ipv4.srcAddr);
  set_metadata(ipv4_metadata.lkp_ipv4_da, ipv4.dstAddr);
  set_metadata(l3_metadata.lkp_ip_proto, ipv4.protocol);
  set_metadata(l3_metadata.lkp_ip_ttl, ipv4.ttl);

  return select(latest.fragOffset, latest.ihl, latest.protocol) {
    0x0000501 : parse_icmp;
    0x0000506 : parse_tcp;
    0x0000511 : parse_udp;
    default   : ingress;
  }
}
```

Metadata can be initialized by the parser

Fields are joined for a match

Deparsing (Serializing packet headers)

- **Fundamental assumption of P4**

- The device must be able to parse any packet it can produce

- **Consequence**

- Packet headers can be reassembled using the parser definition

- When the device only need to insert a header but shouldn't actually parse it

- **Example:** insert my_header after udp

- ```
parser parse_udp {
 extract(udp);
 return select(latest.dst_port) {
 0x0 mask 0x00 : ingress;
 default : parse_my_header;
 }
}
```

Ingress parser will  
always transit to  
ingress

Parser tree has a  
branch to my\_header  
for deparsing

# P4 Language Components

---

- Data declarations
- Parser Programming
- **Packet Flow Programming**
  - Actions
    - Primitive and compound actions
    - Counters, Meters, Registers
  - Tables
    - Match keys
    - Attributes
  - Control Functions

# Actions

---

- **Primitive actions**

- no\_op, drop
- modify\_field, modify\_field\_with\_hash\_based\_offset
- add, add\_to\_field
- add\_header, remove\_header, copy\_header
- push/pop (a header)
- count, execute\_meter
- generate\_digest
- truncate
- resubmit, recirculate, clone{ \_i2i, \_e2i, \_i2e, \_e2e }

- **Compound actions**

```
action route_ipv4(dst_port, dst_mac, src_mac, vid) {
 modify_field(standard_metadata.egress_spec, dst_port);
 modify_field(ethernet.dst_addr, dst_mac);
 modify_field(ethernet.src_addr, src_mac);
 modify_field(vlan_tag.vid, vid);
 add_to_field(ipv4.ttl, -1);
}
```

# Arithmetic and Logical Primitives

---

- **The current standard (v1.0.2)**

- Primitive actions

- Standard: `add()`, `add_to_field()`
    - Additional: `subtract()`, `subtract_from_field()`, `bit_and()`, `bit_or()`, `bit_xor()`, `shift_left()`, `shift_right()`, ...
    - `add_to_field(ipv4.ttl, -1)`

- Partial support for expressions exists in some compilers

- **Developing standard (p4-16)**

- Expressions with `+`, `-`, `&`, `|`, `^`, `~`, `<<`, `>>`, etc.

- `modify_field(ipv4.ttl, ipv4.ttl - 1)`

- Specific targets might restrict expression complexity

# Action Execution Semantics

- All actions within a compound action are assumed to be executed sequentially

```
action parallel_test() {
 modify_field(hdr.fieldA, 1);
 modify_field(hdr.fieldB, hdr.fieldA);
}
```

|        | Sequential Semantics | Parallel Semantics   |
|--------|----------------------|----------------------|
| fieldA | 1                    | 1                    |
| fieldB | 1                    | fieldA before action |

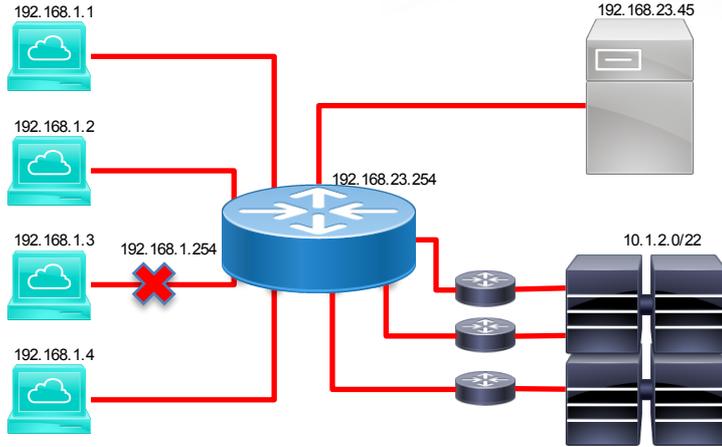
- **This is an important specification change**
  - Up to version 1.0.2 action execution was parallel
  - After 1.0.2 action execution is sequential
- **The maximum number of steps supported for a compound action is target-dependent**

# Match-Action Tables

---

- **The most fundamental units of the Match-Action Pipeline**
- **P4 defines**
  - What to match on and match type
  - A list of *possible* actions
  - Additional attributes
    - Size
- **In runtime, each table contains one or more entries (rows)**
- **An entry contains:**
  - A specific key to match on
  - A **single** action
    - to be executed when a packet matches the entry
  - (Optional) action data

# Example: IPv4 Processing



| Key            | Action         | Action Data                    |
|----------------|----------------|--------------------------------|
| 192.168.1.1    | I3_switch      | port= mac_da= mac_sa= vlan=    |
| 192.168.1.2    | I3_switch      | port= mac_da= mac_sa= vlan=... |
| 192.168.1.3    | I3_drop        |                                |
| 192.168.1.254  | I3_I2_switch   | port=                          |
|                |                |                                |
| 192.168.1.0/24 | I3_I2_switch   | port=                          |
| 10.1.2.0/22    | I3_switch_ecmp | ecmp_group=                    |

- **P4 Program**

- Defines the format of the table
  - Key Fields
  - Actions
  - Action Data

- **Control Plane (IP stack, Routing protocols)**

- Populates table entries with specific information
  - Based on the configuration
  - Based on automatic discovery
  - Based on protocol calculations

- **Data Plane (populated table)**

- Performs the lookup
- Executes the chosen action

# Defining Actions

```
action l3_switch(port, mac_da, mac_sa, vlan) {
 modify_field(metadata.egress_spec, port);
 modify_field(ethernet.dstAddr, mac_da);
 modify_field(ethernet.srcAddr, mac_sa);
 modify_field(vlan_tag[0].vlanid, vlan);
 modify_field(ipv4.ttl, ipv4.ttl - 1);
}

action l3_l2_switch(port) {
 modify_field(metadata.egress_spec, port);
}

action l3_drop() {
 drop();
}

action l3_switch_nexthop(nexthop_index) {
 modify_field(l3_metadata.nexthop, nexthop_index);
 modify_field(l3_metadata.nexthop_type, NEXTHOP_TYPE_SIMPLE);
}

action l3_switch_ecmp(ecmp_group) {
 modify_field(l3_metadata.nexthop, ecmp_group);
 modify_field(l3_metadata.nexthop_type, NEXTHOP_TYPE_ECMP);
}
```

# Match-Action Table (Exact Match)

**Example:** A typical L3 (IPv4) Host table

```
table ipv4_host {
 reads {
 ingress_metadata.vrf : exact;
 ipv4.dstAddr : exact;
 }
 actions {
 l3_switch;
 l3_l2_switch;
 l3_switch_nexthop;
 l3_switch_ecmp;
 l3_drop;
 }
 size : HOST_TABLE_SIZE;
}
```



These are the only possible actions. Each particular entry can have only ONE of them.

| vrf | ipv4.dstAddr | action         | data                     |
|-----|--------------|----------------|--------------------------|
| 1   | 192.168.1.10 | l3_switch      | port_id= mac_da= mac_sa= |
| 100 | 192.168.1.10 | l3_l2_switch   | port_id=<CPU>            |
| 1   | 192.168.1.3  | l3_drop        |                          |
| 5   | 10.10.1.1    | l3_switch_ecmp | ecmp_group=127           |

# Match-Action Table (Longest Prefix Match)

**Example:** A typical L3 (IPv4) Routing table

```
table ipv4_lpm {
 reads {
 ingress_metadata.vrf : exact;
 ipv4.dstAddr : lpm;
 }
 actions {
 l3_l2_switch;
 l3_multicast;
 l3_nextHop;
 l3_ecmp;
 l3_drop;
 }
 size : 65536;
}
```

Different fields can use different match types

Prefix also serves as a priority indicator

| vrf | ipv4.dstAddr / prefix | action            | data              |
|-----|-----------------------|-------------------|-------------------|
| 1   | 192.168.1.0 / 24      | l3_l2_switch      | port_id=64        |
| 10  | 10.0.16.0 / 22        | l3_ecmp           | ecmp_index=12     |
| 1   | 192.168.0.0 / 16      | l3_switch_nextHop | nextHop_index=451 |
| 1   | 0.0.0.0 / 0           | l3_switch_nextHop | nextHop_index=1   |

# Match-Action Table (Ternary Match)

**Example:** A typical L3 (IPv4) Routing table

```
table ipv4_lpm {
 reads {
 ingress_metadata.vrf : ternary;
 ipv4.dstAddr : ternary;
 }
 actions {
 l3_l2_switch;
 l3_multicast;
 l3_nexthop;
 l3_ecmp;
 l3_drop;
 }
 size : 65536;
}
```

Ternary tables require an explicit specification of entry priority

| Prio | vrf / mask   | ipv4.dstAddr / mask           | action            | data             |
|------|--------------|-------------------------------|-------------------|------------------|
| 100  | 0x001/0xFFFF | 192.168.1.5 / 255.255.255.255 | l3_switch_nexthop | nexthop_index=10 |
| 10   | 0x000/0x000  | 192.168.2.0/255.255.255.0     | l3_switch_ecmp    | ecmp_index=25    |
| 10   | 0x000/0x000  | 192.168.3.0/255.255.255.0     | l3_switch_nexthop | nexthop_index=31 |
| 5    | 0x000/0x000  | 0.0.0.0/0.0.0.0               | l3_l2_switch      | port_id=64       |

# Match Types

---

- **Exact**
  - port\_index : exact
- **Ternary**
  - ethernet.srcAddr : ternary
- **LPM (special kind of ternary match)**
  - ipv4.dstAddr : lpm
- **Range**
  - udp.dstPort : range
- **Valid**
  - vlan\_tag[0] : valid

# Table Miss

---

- **Each table can have a Default Action**

- Chosen by the Control Path at runtime from the list of table Actions
  - P4 Program does not have an indication which action (and which action data) will be the default

- **When no matching entries are found**

- Default Action with the default action data is executed **if** it has been set by the control path
- If no Default Action has been specified, it is **no\_op()**

# Stateful Objects

---

- **Counters, Meters, Registers**

# What are stateful objects

---

- **Stateful objects keep their state between packets**
  - Metadata and packet headers are **stateless**
    - They are re-initialized for each packet
  - Counters, Meters and Registers are **stateful**
    - Counters are incremented with each packet
    - Meters keep their bucket state
    - Registers store arbitrary data

# Direct Counters

## A counter per table entry

```
counter ip_acl_stats {
 type : packets_and_bytes;
 direct : ip_acl;
}
```

```
table ip_acl {
 reads {
 ipv4_metadata.lkp_ipv4_sa : ternary;
 ipv4_metadata.lkp_ipv4_da : ternary;
 l3_metadata.lkp_ip_proto : ternary;
 l3_metadata.lkp_l4_sport : ternary;
 l3_metadata.lkp_l4_dport : ternary;
 }
 actions {
 nop;
 acl_log;
 acl_deny;
 acl_permit;
 acl_mirror;
 acl_redirect_nexthop;
 acl_redirect_ecmp;
 }
 size : INGRESS_IP_ACL_TABLE_SIZE;
}
```

| Match Fields         | Action Sel        | Action Data  | counter ip_acl_stats Counter |
|----------------------|-------------------|--------------|------------------------------|
| ABCD_xxxx_0123       | acl_deny          |              | counter A                    |
| <b>matched entry</b> | <b>acl_permit</b> | <b>8b 8b</b> | <b>pkt/byte counts</b>       |
|                      |                   |              |                              |
| BA8E_F007_xxxx       | nop               |              | counter Z                    |

# Indirect Counters

## Flexibly linked counters

```
counter ingress_bd_stats {
 type : packets_and_bytes;
 instance_count : BD_STATS_TABLE_SIZE;
}

action set_bd(bd, bd_stat_index) {
 modify_field(12_metadata.bd, bd);
 count(ingress_bd_stats, bd_stat_index);
}

table port_vlan {
 reads {
 ingress_metadata.ingress_port : exact;
 vlan_tag[0] : valid;
 vlan_tag[0].vlan_id : exact;
 }
 actions {
 set_bd;
 }
}
```

Different VLANs (BDs) can share the same counter

Other tables can also reference these counters

table port\_vlan

| Match Fields         | Action Sel    | Action Data               |
|----------------------|---------------|---------------------------|
| ABCD_0123            | set_bd        | bd bd_stat_index          |
|                      | set_bd        | bd bd_stat_index          |
| <b>matched entry</b> | <b>set_bd</b> | <b>bd bd_stat_index A</b> |
|                      | set_bd        | bd bd_stat_index          |
|                      | set_bd        | bd bd_stat_index          |
|                      | set_bd        | bd bd_stat_index A        |
| BA8E_F007            | set_bd        | bd bd_stat_index          |

counter ingress\_bd\_stats

|                        |
|------------------------|
|                        |
|                        |
| <b>pkt/byte counts</b> |
|                        |
|                        |

# Meters

- Declaration is similar to counters

- Action: `execute_meter(meter_array, meter_index, color_destination)`

```
/* Direct Meter */
meter acl_meter {
 type: packets;
 direct: ip_acl;
 result: metadata.color;
}
```

Meters calculate packet color and deposit it into the specified field

```
/* Indirect Meter Array */
meter bd_meter {
 type: bytes;
 instance_count: 1000;
}
```

```
action do_bd_meter(meter_index) {
 execute_meter(bd_meter, meter_index, metadata.color);
}
```

Color Coding:

0 – Green  
1 – Yellow  
2 -- Red

# Registers

- **Declaration is similar to indirect counters**

- Actions

- `register_read(register_array, register_index, destination_field)`
- `register_write(register_array, register_index, value)`

```
/* Register Array (Indirect) */
register last_syn {
 width: 32;
 static: flow_table;
 instance_count: 1024;
}

action get_flow_age(flow_index) {
 register_read(last_syn, flow_index, metadata.flow_start_time);
 modify_field(metadata.flow_age,
 metadata.flow_start_time - metadata.ingress_global_stamp);
}

action start_new_flow(flow_index) {
 register_write(last_syn, flow_index, metadata.ingress_global_timestamp);
}
```

# Control Flow Functions

---

- **Primitives**

- Perform a table lookup: **apply**
- **if/else** statement
- **apply** with the case clause

- **Sequential Execution Semantics**

- Compiler is doing parallelization automatically

- **Standard control functions**

- `ingress()` – Ingress Pipeline processing
- `egress()` – Egress Pipeline processing

- **User-defined control functions**

# Standard Control Functions

---

- **ingress()** control function starts processing
  - remember “return ingress;” statement in the parser functions
- **egress()** control function is called implicitly from the Packet Replication Engine

```
control ingress {
 apply(port_vlan_mapping);
 apply(smac);
 apply(dmac);
}
```

```
control egress {
 apply(vlan_tag_removal);
}
```

# User-Defined Control Functions

- **Help improve code readability**
  - No specific performance advantages: the code is flattened by the compiler
- **No parameters are accepted**

```
control assign_vlan {
 apply(subnet_vlan);
 apply(mac_vlan);
 apply(protocol_vlan);
 apply(port_vlan);
 apply(resolve_vlan);
}

control ingress {
 . . .
 if (!valid(vlan_tag[0]) {
 assign_vlan();
 }
 . . .
}
```

# If/Else Branching

## Example: Separate Ipv4 and IPv6 Processing Paths

```
if ((l3_metadata.lkp_ip_type == IPTYPE_IPV4) and (ipv4_metadata.ipv4_unicast_enabled == TRUE)) {
 process_ipv4_racl();
 process_nat();
 process_ipv4_urpf();
 process_ipv4_fib();
}
else {
 if ((l3_metadata.lkp_ip_type == IPTYPE_IPV6) and (ipv6_metadata.ipv6_unicast_enabled == TRUE)) {
 process_ipv6_racl();
 process_ipv6_urpf();
 process_ipv6_fib();
 }
}
```

# Action Branching

**Example:** Use per-router-mac decapsulation

```
table router_mac {
 reads {
 l2_metadata.lkup_dst_mac : ternary;
 l2_metadata.bd : ternary;
 ingress_metadata.src_port: ternary;
 }
 actions {
 nop;
 enable_ipv4_lookup;
 enable_ipv6_lookup;
 enable_mpls_decap;
 enable_mim_decap;
 }
}

control process_router_mac_lookup {
 apply(router_mac) {
 enable_ipv4_lookup { process_ipv4_fib(); }
 enable_ipv6_lookup { process_ipv6_fib(); }
 enable_mpls_decap { process_mpls_label_lookup(); }
 /* etc. */
 }
}
```

# Miss branching

```
action on_miss() {}

table ipv4_fib {
 reads {
 . . .
 }
 actions {
 13_switch;
 13_l2_switch;
 13_switch_nexthop;
 13_switch_ecmp;
 on_miss;
 }
}

control process ipv4_fib {
 apply(ipv4_fib) {
 on_miss {
 apply(ipv4_fib_lpm);
 }
 }
}
```

We choose to use only this action as the default

# Executing actions in the control flow

```
control process_counters {
 if (my_meta.drop_packet == 0) {
 count(bd_counter,
 metadata.bd_counter_index);
 count(vrf_counter,
 metadata.vrf_counter_index);
 }
}
```

```
action update_counters () {
 count(bd_counter,
 metadata.bd_counter_index);
 count(vrf_counter,
 metadata.vrf_counter_index);
}
```

```
table do_process_counters {
 actions {
 update_counters;
 }
}
```

This action must be set as a default for this table by the control plane

```
control process_counters {
 if (my_meta.drop_packet == 0) {
 apply(do_process_counters);
 }
}
```

- **Actions cannot be directly referenced in the control flow functions**
  - Instead, they need to be “wrapped” into tables
- **Tables without keys can be used to implement unconditional execution**
  - They always miss and hence the desired action needs to be set as a default

# Advanced Concepts

---

- **Action Profiles**
- **Packet Digests**
- **Packet Resubmit/Recirculation**
- **Packet Cloning**

# Action Profiles

---

- **Separate table match entries from actions and action data**
- **Allow multiple entries to share same action data**
  - Saves space
  - Allows quick update of multiple entries
- **Allow multiple actions/action\_data per entry**
  - This is called “dynamic action selection”
  - Used to implement LAG or ECMP
- **Can be more efficient compared to explicit implementation**

# Action Profiles

Actions can be complex

```
action set_bd(bd, vrf, rmac_group,
 ipv4_unicast_enabled, ipv6_unicast_enabled,
 ipv4_urpf_mode, ipv6_urpf_mode,
 igmp_snooping_enabled, mld_snooping_enabled,
 bd_label, stp_group, stats_idx,
 exclusion_id)
{
 modify_field(l3_metadata.vrf, vrf);
 modify_field(ipv4_metadata.ipv4_unicast_enabled,
 modify_field(ipv6_metadata.ipv6_unicast_enabled,
 modify_field(ipv4_metadata.ipv4_urpf_mode,
 modify_field(ipv6_metadata.ipv6_urpf_mode,
 modify_field(l3_metadata.rmac_group,
 modify_field(acl_metadata.bd_label,
 modify_field(ingress_metadata.bd,
 modify_field(ingress_metadata.outer_bd,
 modify_field(l2_metadata.stp_group,
 modify_field(l2_metadata.bd_stats_idx,
 modify_field(multicast_metadata.igmp_snooping_enabled,
 modify_field(multicast_metadata.mld_snooping_enabled,
 modify_field(ig_intr_md_for_tm.level1_exclusion_id,
}
```

60-70 bits for the  
parameters







# Using the profiles for LAG and ECMP

```
action_selector ecmp_selector {
 selection_key : ecmp_hash;
}

action_profile ecmp_action_profile {
 actions {
 nop;
 set_ecmp_nexthop_details;
 }
 size : ECMP_SELECT_TABLE_SIZE;
 dynamic_action_selection : ecmp_selector;
}

table ecmp_group {
 reads {
 13_metadata.nexthop_index : exact;
 }
 action_profile : ecmp_action_profile;
 size : ECMP_GROUP_TABLE_SIZE;
}
```

Chooses a particular entry within a group

Chooses a GROUP of profile entries

# Using the profiles for LAG and ECMP

```
action_selector ecmp_selector {
 selection_key : ecmp_hash;
}

action_profile ecmp_action_profile {
 actions {
 nop;
 set_ecmp_nexthop_details;
 }
 size : ECMP_SELECT_TABLE_SIZE;
 dynamic_action_selection : ecmp_selector;
}

table ecmp_group {
 reads {
 l3_metadata.nexthop_index : exact;
 }
 action_profile : ecmp_action_profile;
 size : ECMP_GROUP_TABLE_SIZE;
}
```

```
field_list l3_hash_fields {
 ipv4_metadata.lkp_ipv4_sa;
 ipv4_metadata.lkp_ipv4_da;
 l3_metadata.lkp_ip_proto;
 l3_metadata.lkp_l4_sport;
 l3_metadata.lkp_l4_dport;
}

field_list_calculation ecmp_hash {
 input {
 l3_hash_fields;
 }
 algorithm : crc16;
 output_width : ECMP_BIT_WIDTH;
}
```

# Packet Digests – Notify control plane of data plane events

---

- **Action**

- generate\_digest(receiver, field\_list)

- **Sends the specified metadata to a target-specific receiver**

- Software-based

- Hardware-Based

# Implementing MAC Learning

## Source MAC Lookup

```
action smac_hit(src_port, static) {
 modify_field(l2_metadata.src_port, src_port);
 modify_field(l2_metadata.is_static, static);
 modify_field(l2_metadata.src_move,
 standard_metadata.ingress_port ^ src_port);
}

action smac_miss() {
 modify_field(l2_metadata.src_miss, TRUE);
}

table smac {
 reads {
 vlan_tag[0].vid : exact;
 ethernet.srcAddr: exact;
 }
 actions {
 smac_hit;
 smac_miss;
 }
}
```

## Learn Notification Generation

```
field_list l2_learn_digest {
 ethernet.srcAddr,
 vlan_tag[0].vid,
 standard_metadata.ingress_port,
 l2_metadata.src_port,
 l2_metadata.src_move
}

action send_learn_notification() {
 generate_digest(L2_RECV, l2_learn_digest);
}

table learn_notification {
 reads {
 l2_metadata.src_move : exact;
 l2_metadata.src_miss : exact;
 l2_metadata.is_static: exact;
 }
 actions {
 send_learn_notification;
 }
}

control process_l2_learning {
 apply(smac);
 apply(learn_notification);
}
```

Target-specific number. Can be SW or HW receiver

Other fields can influence sending of learn notifications

# Resubmit and Recirculate

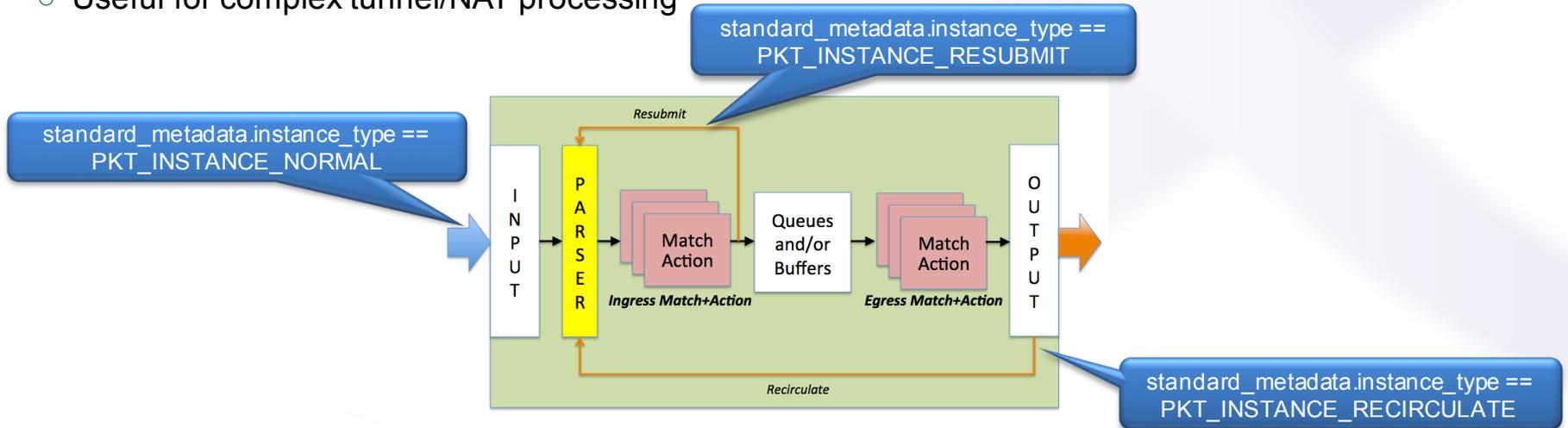
- Implementing complex processing

- **resubmit(field\_list)**

- **Unmodified** packet is returned back to the parser within ingress pipeline, with additional metadata
- Useful for complex parsing

- **recirculate(field\_list)**

- **Fully modified packet** is returned back to the ingress parser after completing all the processing
- Useful for complex tunnel/NAT processing



# MPLS Processing (Simple)

```
parser parse_mpls {
 extract(mpls[next]);
 return select(latest.bos) {
 0 : parse_mpls;
 1 : parse_mpls_bos;
 default : ingress;
 }
}

parser parse_mpls_bos {
 return select(current(0, 4)) {
 0x4 : parse_inner_ipv4;
 0x6 : parse_inner_ipv6;
 default: parse_inner_ethernet;
 }
}
```

- **This parser is a very effective hack**
  - MPLS does not formally carry any indication of the payload type
- **The right way to do parsing is by using label lookup**

# MPLS Processing (Full)

## Parser

```
#define IPV4 0
#define IPV6 1
#define ETHERNET 2

header_type my_metadata {
 fields {
 mpls_payload : 2;
 }
}

parser parse_mpls_bos {
 return select(
 standard_metadata.resubmit_flag,
 my_metadata.mpls_payload,
 current(0, 4)) {
 0x04 : parse_inner_ipv4;
 0x06 : parse_inner_ipv6;
 0x40 mask 0x70 : parse_inner_ipv4
 0x50 mask 0x70 : parse_inner_ipv6;
 0x60 mask 0x70 : parse_inner_ethernet;
 default: parse_inner_ethernet;
 }
}
```

## Control Flow

```
action set_mpls_payload_type(mpls_payload) {
 modify_field(my_metadata.mpls_payload, mpls_payload);
}

table mpls {
 reads {
 mpls[0].label : exact;
 mpls[1].label : exact;
 }
 actions {
 set_mpls_payload_type;
 }
}

control ingress {
 . . .
 apply(mpls);
 /* If our guess was wrong... */
 if (my_metadata.mpls_payload == ETHERNET and !valid(inner_ethernet)) {
 apply(mpls_resubmit);
 }
 . . .
}
```

# MPLS Processing (Resubmit)

```
field_list mpls_resubmit_list {
 my_metadata.mpls_payload
}

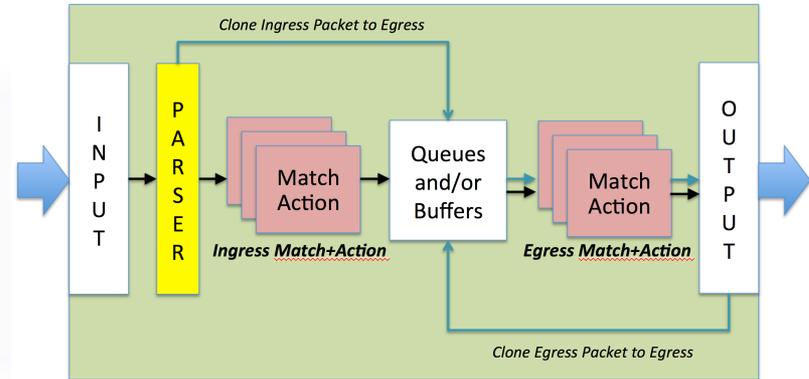
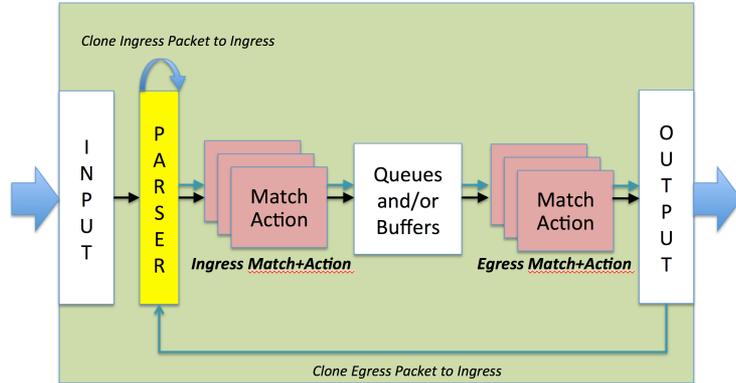
action do_mpls_resubmit() {
 resubmit(mpls_resubmit_list);
}

table mpls_resubmit {
 actions {
 do_mpls_resubmit;
 }
}
```

- **The packet still goes through the rest of the ingress pipeline**
  - Use the if() statement around the rest of the code to avoid further processing
- **All metadata computed at the first pass will be lost**
  - Put all the fields that will be needed in the second pass into the resubmit field\_list

# Cloning

- A new copy of a packet is created
- The original and the clone are processed independently
- Two cloning sources
  - Original packet (before ingress)
  - Fully processed (after egress)
- Two cloning destinations
  - Beginning of the ingress pipeline
  - Packet Queuing/Replication Engine
    - Before the beginning of the egress pipeline



| Dest \ Source    | Packet before Ingress | Packet After Egress |
|------------------|-----------------------|---------------------|
| Ingress Pipeline | clone_i2i()           | clone_e2i()         |
| Egress Pipeline  | clone_i2e()           | clone_e2e()         |

`clone_x2y(clone_spec, field_list)`

# Implementing Negative Mirroring (aka mirror on drop)

## Creating a mirrored copy

```
field_list mirror_info {
 i2e_metadata.mirror_session_id;
 i2e_metadata.mirror_drop_flags;
}

action do_negative_mirror(session_id) {
 modify_field(i2e_metadata.mirror_session_id, session_id);
 modify_field(i2e_metadata.mirror_drop_flags,
 ingress_metadata.drop_flags);
 clone_i2e(session_id, mirror_info);
}

table negative_mirror {
 actions {
 do_negative_mirror;
 }
}

control ingress {
 . . .
 if (ingress_metadata.drop_flags != 0) {
 apply(negative_mirror);
 }
}
```

## Processing the mirrored copy

```
action set_mirror_nhop(nhop_idx) {
 modify_field(l3_metadata.nexthop_index, nhop_idx);
}

action set_mirror_bd(bd) {
 modify_field(egress_metadata.bd, bd);
}

table mirror {
 reads {
 i2e_metadata.mirror_session_id : exact;
 }
 actions {
 nop;
 set_mirror_nhop;
 set_mirror_bd;
 }
 size : MIRROR_SESSIONS_TABLE_SIZE;
}

control egress {
 if (standard_metadata.instance_type == PKT_INSTANCE_TYPE_INGRESS_CLONE) {
 apply(mirror);
 }
 else {
 . . .
 }
}
```

# P4 Compiler Overview

# Modular Compiler Overview

---

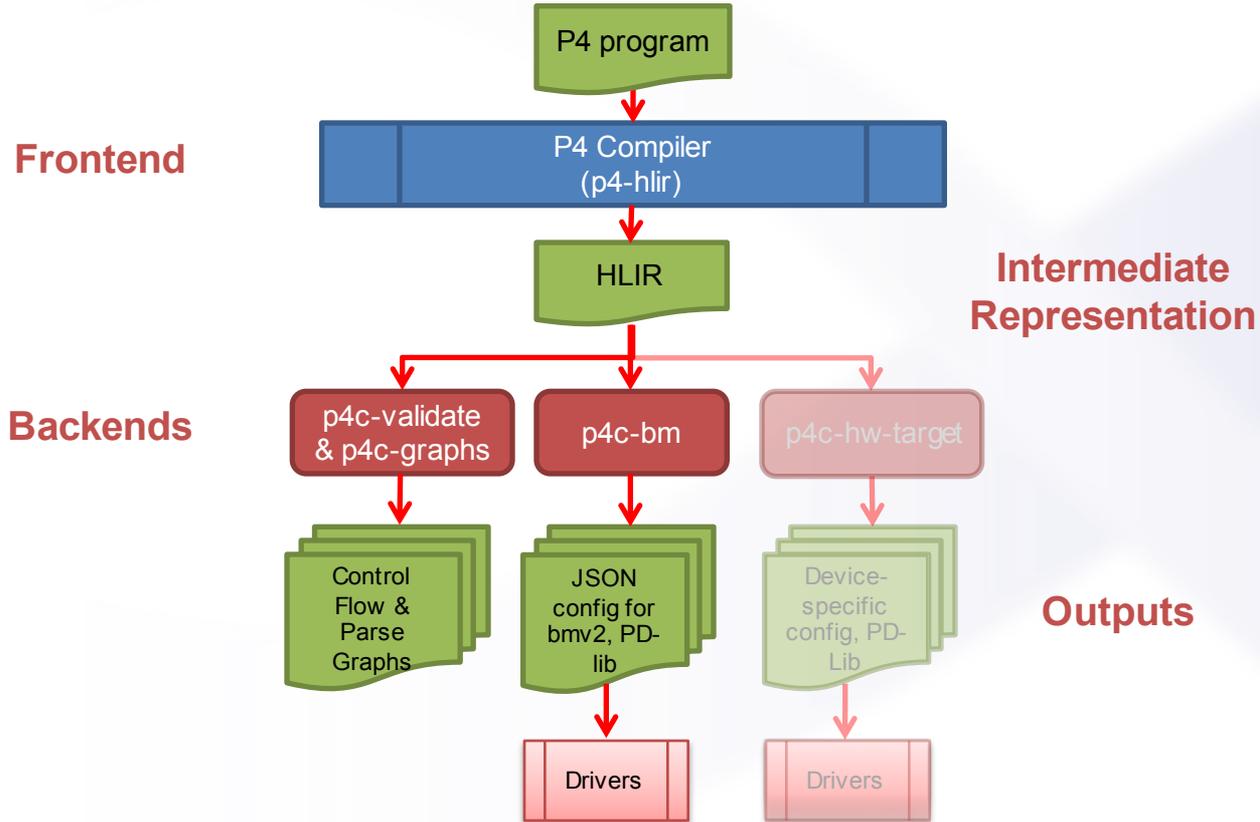
- **Single Front-End (p4-hlir)**

- Translates P4 code into High-Level Intermediate Representation (HLIR)
  - Similar to AST (Abstract Syntax Trees)
  - Currently represented as a hierarchy of Python objects
  - Frees backend developers from the burden of syntax analysis and target-independent semantic checks
  - HLIR documentation is supplied with the frontend code

- **Multiple backends**

- Code generators for various targets
  - Software Switch Model (p4c-bm)
  - Network Interface Cards
  - Packet Processors / NPUs
  - FPGAs, GPUs, ASICs
- Validators and graph generators
- Run-time API generators

# P4 Modular Compiler



# Dependency Analysis

# Types of dependencies

---

- **Dependencies are inferred from target-independent P4 program analysis**
- **Independent tables**
- **Match Dependency**
- **Action Dependency**
- **Successor Dependency**
- **Reverse Read Dependency**

# Independent Tables

```
action ing_drop() {
 modify_field(ing_metadata.drop, 1);
}

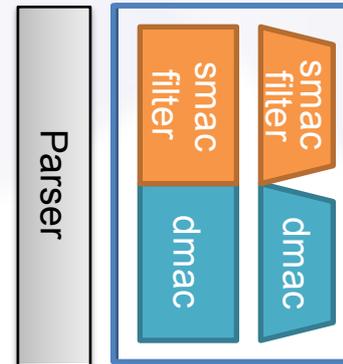
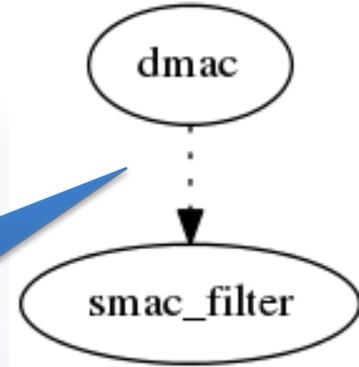
action set_egress_port(egress_port) {
 modify_field(ing_metadata.egress_spec, egress_port);
}

table dmac {
 reads {
 ethernet.dstAddr : exact;
 }
 actions {
 nop;
 set_egress_port;
 }
}

table smac_filter {
 reads {
 ethernet.srcAddr : exact;
 }
 actions {
 nop;
 ing_drop;
 }
}

control ingress {
 apply(dmac);
 apply(smac_filter);
}
```

Tables are independent: both matching and action execution can be done in parallel



# Action Dependency

```
action ing_drop() {
 modify_field(ing_metadata.drop, 1);
}

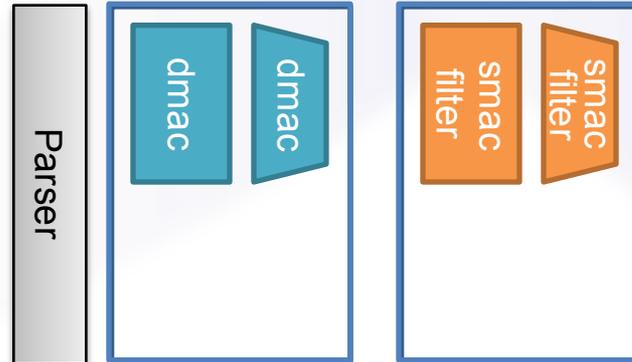
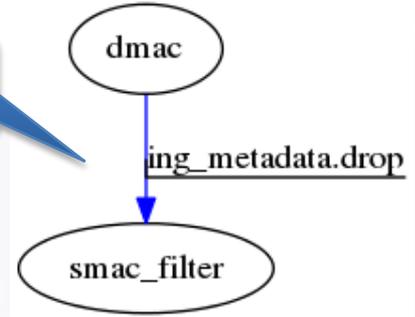
action set_egress_port(egress_port) {
 modify_field(ing_metadata.egress_spec, egress_port);
}

table dmac {
 reads {
 ethernet.dstAddr : exact;
 }
 actions {
 nop;
 ing_drop;
 set_egress_port;
 }
}

table smac_filter {
 reads {
 ethernet.srcAddr : exact;
 }
 actions {
 nop;
 ing_drop;
 }
}

control ingress {
 apply(dmac);
 apply(smac_filter);
}
```

Tables act on the same field and therefore must be placed in separate stages



# Match Dependency

```
action set_bd(bd) {
 modify_field(ing_metadata.bd, bd);
}

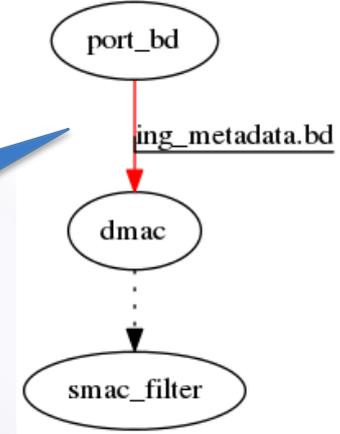
table port_bd {
 reads {
 ing_metadata.ingress_port : exact;
 }
 actions {
 set_bd;
 }
}

table dmac {
 reads {
 ethernet.dstAddr : exact;
 ing_metadata.bd : exact;
 }
 actions {
 nop;
 set_egress_port;
 }
}

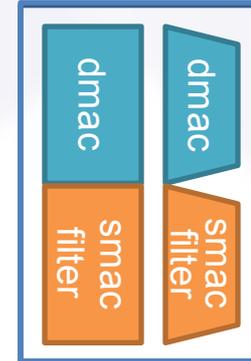
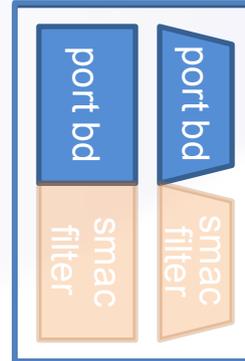
table smac_filter {
 reads {
 ethernet.srcAddr : exact;
 }
 actions {
 nop;
 ing_drop;
 }
}

control ingress {
 apply(port_bd);
 apply(dmac);
 apply(smac_filter);
}
```

The second table matches on the field, modified by the first.



Parser



# Successor Dependency

```
action set_bd(bd) {
 modify_field(ing_metadata.bd, bd);
}

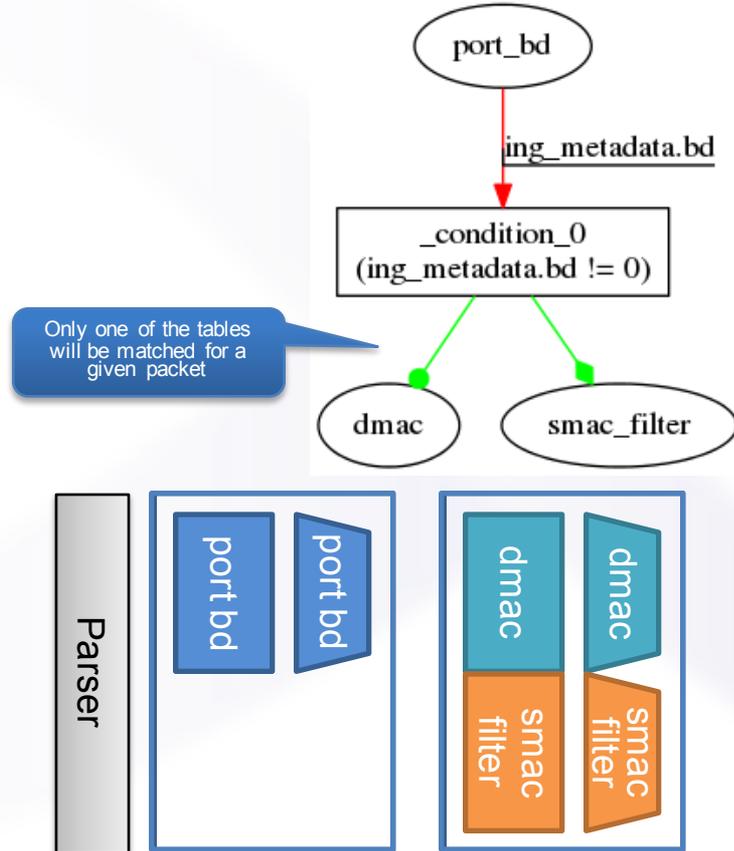
table port_bd {
 reads {
 ing_metadata.ingress_port : exact;
 }
 actions {
 set_bd;
 }
}

table dmac {
 reads {
 ethernet.dstAddr : exact;
 ing_metadata.bd : exact;
 }
 actions {
 nop;
 ing_drop;
 set_egress_port;
 }
}

table smac_filter {
 reads {
 ethernet.srcAddr : exact;
 }
 actions {
 nop;
 ing_drop;
 }
}

control ingress {
 apply(port_bd);

 if (ing_metadata.bd != 0) {
 apply(dmac);
 } else {
 apply(smac_filter);
 }
}
```



# Reverse Read Dependency

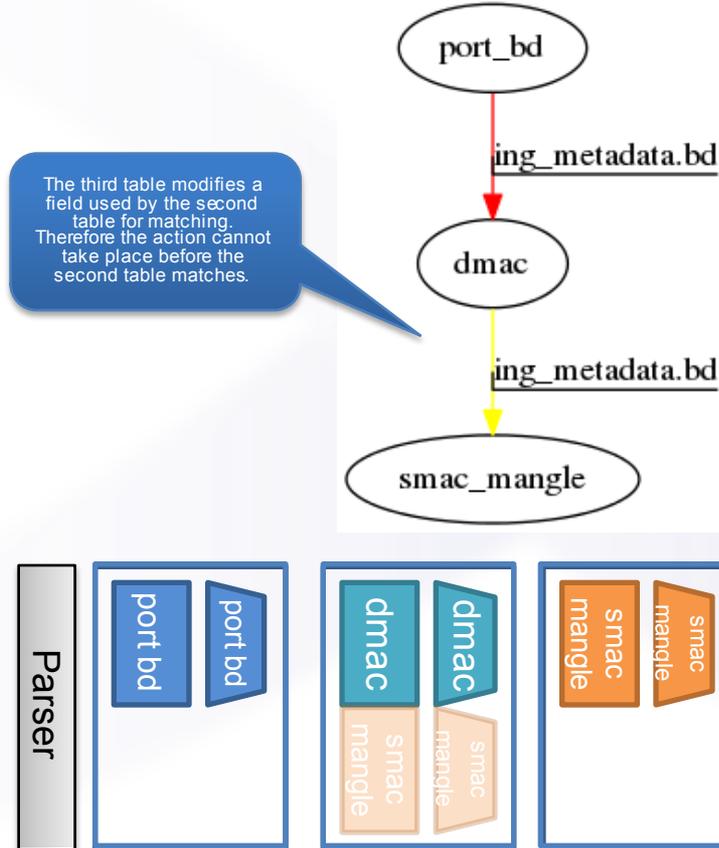
```
action set_bd(bd) {
 modify_field(ing_metadata.bd, bd);
}

table port_bd {
 reads {
 ing_metadata.ingress_port : exact;
 }
 actions {
 set_bd;
 }
}

table dmac {
 reads {
 ethernet.dstAddr : exact;
 ing_metadata.bd : exact;
 }
 actions {
 nop;
 ing_drop;
 set_egress_port;
 }
}

table smac_mangle {
 reads {
 ethernet.srcAddr : exact;
 }
 actions {
 nop;
 set_bd;
 }
}

control ingress {
 apply(port_bd);
 apply(dmac);
 apply(smac_mangle);
}
```

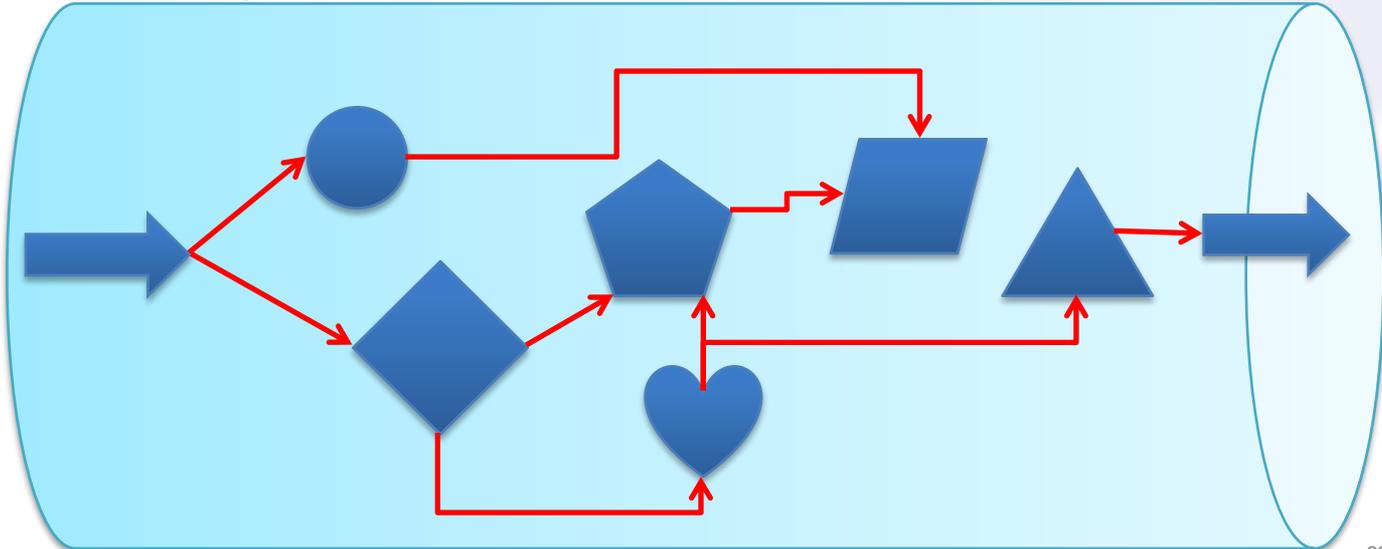


# Automatic API Generation

---

# Network Device API Basics

- **Object Definitions (Schema)**
  - Reflects the object properties and methods
- **Object Relationships (Behavior)**
  - The quality of the API is directly dependent on how well the object relationships are specified



# P4 is an Ideal Base for a Network APIs

---

- **Clearly defined objects**

- Tables
- Counters
- Meters
- Registers

- **Unambiguously defined relationships**

- Control Flow Functions

- **Idea:**

- Each of fundamental P4 objects has a “natural” schema



# Example API. Match & Action Specs

## myprog.p4

```
action a1(p11, p12) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
}
```

## pd\_myprog.h

```
typedef struct p4_pd_myprog_a1_action_spec {
 <type> p11;
 <type> p12;
} p4_pd_myprog_a1_action_spec_t;

typedef struct p4_pd_myprog_a2_action_spec {
 <type> p21;
 <type> p22;
 <type> p23;
} p4_pd_myprog_a2_action_spec_t;

typedef struct p4_pd_myprog_a3_action_spec {
} p4_pd_myprog_a3_action_spec_t;

typedef struct p4_pd_myprog_t1_match_spec {
 <type> meta_f1;
 <type> meta_f2;
 <type> meta_f2_mask;
 uint8_t h1_valid;
} p4_pd_myprog_t1_match_spec_t;
```

**exact:** f  
**ternary:** f and f\_mask  
**lpm:** f and f\_prefix\_len  
**valid:** f\_valid  
**range:** f\_min and f\_max

# Example API. Entry Add

## myprog.p4

```
action a1(p11, p12) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
 actions {
 a1;
 a2;
 a3;
 }
}
```

## pd\_myprog.h

```
p4_pd_status_t p4_pd_myprog_t1_entry_add_with_a1(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_priority_t priority,
 const p4_pd_myprog_t1_match_spec_t *match_spec,
 const p4_pd_myprog_a1_action_spec_t *action_spec,
 p4_pd_entry_handle_t *entry_hdl);

p4_pd_status_t p4_pd_myprog_t1_entry_add_with_a2(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_priority_t priority,
 const p4_pd_myprog_t1_match_spec_t *match_spec,
 const p4_pd_myprog_a2_action_spec_t *action_spec,
 p4_pd_entry_handle_t *entry_hdl);

p4_pd_status_t p4_pd_myprog_t1_entry_add_with_a3(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_priority_t priority,
 const p4_pd_myprog_t1_match_spec_t *match_spec,
 const p4_pd_myprog_a3_action_spec_t *action_spec,
 p4_pd_entry_handle_t *entry_hdl);
```

# Example API. Entry Modify

## myprog.p4

```
action a1(p11, p2) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
 actions {
 a1;
 a2;
 a3;
 }
}
```

## pd\_myprog.h

```
p4_pd_status_t p4_pd_myprog_t1_entry_modify_with_a1(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl,
 const p4_pd_myprog_a1_action_spec_t *action_spec);

p4_pd_status_t p4_pd_myprog_t1_entry_modify_with_a2(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl,
 const p4_pd_myprog_a2_action_spec_t *action_spec);

p4_pd_status_t p4_pd_myprog_t1_entry_modify_with_a3(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl,
 const p4_pd_myprog_a3_action_spec_t *action_spec);
```

# Example API. Entry Delete and Lookup

## myprog.p4

```
action a1(p11, p12) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
 actions {
 a1;
 a2;
 a3;
 }
}
```

## pd\_myprog.h

```
p4_pd_status_t p4_pd_myprog_t1_entry_delete(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl);

p4_pd_status_t p4_pd_myprog_t1_entry_lookup(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 const p4_pd_myprog_t1_match_spec_t *match_spec,
 p4_pd_entry_handle_t *entry_hdl);
```

# Example API. Entry Get

## myprog.p4

```
action a1(p11, p12) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
 actions {
 a1;
 a2;
 a3;
 }
}
```

## pd\_myprog.h

```
typedef enum {
 P4_PD_MYPROG_ACTION_A1,
 P4_PD_MYPROG_ACTION_A2,
 P4_PD_MYPROG_ACTION_A3,
 ...
 P4_PD_MYPROG_ACTION_COUNT;
} p4_pd_myprog_actions_t;

typedef union {
 p4_pd_myprog_a1_action_spec_t a1;
 p4_pd_myprog_a2_action_spec_t a2;
 . . .
} p4_pd_myprog_action_spec_t;

p4_pd_status_t p4_pd_myprog_t1_entry_get(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl,
 p4_pd_myprog_t1_match_spec_t *match_spec,
 p4_pd_myprog_actions_t *action,
 p4_pd_myprog_action_spec_t *action_spec_t);
```

# Example API. Default Action APIs

## myprog.p4

```
action a1(p11, p2) {...}
action a2(p21, p22, p23) {...}
action a3() {...}

table t1 {
 reads {
 meta.f1 : exact;
 meta.f2 : ternary;
 h1 : valid;
 }
 actions {
 a1;
 a2;
 a3;
 }
}
```

## pd\_myprog.h

```
p4_pd_status_t p4_pd_myprog_t1_set_default_action_a1(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 const p4_pd_myprog_a1_action_spec_t *action_spec);

p4_pd_status_t p4_pd_myprog_t1_set_default_action_a2(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 const p4_pd_myprog_a2_action_spec_t *action_spec);

p4_pd_status_t p4_pd_myprog_t1_set_default_action_a3(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 const p4_pd_myprog_a3_action_spec_t *action_spec);

p4_pd_status_t p4_pd_myprog_t1_clear_default_action(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle);
```

# Counters

---

- **Individual Counter Operations**

- Get
  - (Counter Index or Entry ID) → Value
- Clear
  - (Counter Index or Entry ID) →
- Set (optional)
  - (Counter Index or Entry ID, Value) →

- **Counter Array Operations**

- Width Get
- Array size Get
- Get All
- Clear All

# Example API

## myprog.p4

```
counter c1 {
 type: packets_and_bytes;
 direct: t1;
};

counter c2 {
 type: bytes;
 instance_count: 1000;
}
```

## pd\_myprog.h

```
p4_pd_status_t p4_pd_myprog_c1_get(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 p4_pd_entry_handle_t entry_hdl,
 uint64_t *packets,
 uint64_t *bytes);

p4_pd_status_t p4_pd_myprog_c2_get(
 p4_pd_target_t device_target,
 p4_pd_session_t session_handle,
 uint32_t counter_idx,
 uint64_t *bytes);
```

# Meters

---

- **Individual Meter Operations**

- **Set**

- (Meter Index or EntryID, Committed Rate, Committed Birst, Peak Rate, Peak Birst)
    - Is that the only option?
    - What about different meter types (color-blind/color-aware, single rate?)
      - Are all meters in the array of the same type?
    - Who standardizes the units (bits, bytes, kbits, Mbytes, etc.)?
    - Who standardizes the colors?

- **Get**

- (Meter Index or Entry ID) → (Settings)

# Registers

---

- **Operations**

- Set

- (Register Index or Entry ID, value) →

- Get

- (Register Index or Entry ID) → value

- **C type for the value depends on register definition**

- **Optional Operations**

- Width Get

- Get All

- Set All

# github.com/p4lang

---

- **switch.p4: reference p4 program**
- **BMv2 (Behavioral Model v2): s/w switch runs p4**
- **BMv2 compiler**
- **supporting tools, scripts**
  
- **Project summaries: [link](#)**

# BMv2 Primitives

---

- standard primitives [https://github.com/p4lang/p4-hlr/blob/master/p4\\_hlr/frontend/primitives.json](https://github.com/p4lang/p4-hlr/blob/master/p4_hlr/frontend/primitives.json)
- bmv2 specific primitives [https://github.com/p4lang/p4c-bm/blob/master/p4c\\_bm/primitives.json](https://github.com/p4lang/p4c-bm/blob/master/p4c_bm/primitives.json)

# Intrinsic Metadata, provided by BMv2 switch

- If one defines all these fields, all the simple\_switch features will be supported, so it is recommended to define these fields in every program (to avoid a headache).

```
header_type intrinsic_metadata_t{
 fields{
 ingress_global_timestamp : 48; // ingress timestamp, in microseconds
 mcast_grp : 4; // to be set in the ingress pipeline for multicast
 egress rid : 4; // replication id, available on egress if packet was multicast
 mcast_hash : 16; // unused
 lf_field_list : 32; // set by generate_digest primitive, not to be used directly by P4 programmer
 resubmit_flag : 16; // used internally
 recirculate_flag : 16; // used internally
 }
}
```

```
metadata intrinsic_metadata_t intrinsic_metadata;
```

```
header_type queueing_metadata_t{
 fields{
 enq_timestamp : 48; // in microseconds
 enq_qdepth : 16;
 deq_timedelta : 32;
 deq_qdepth : 16;
 }
}
```

```
metadata queueing_metadata_t queueing_metadata;
```

**Thank you**

